

Addressing heterogeneity in model-based development of cyber-physical systems

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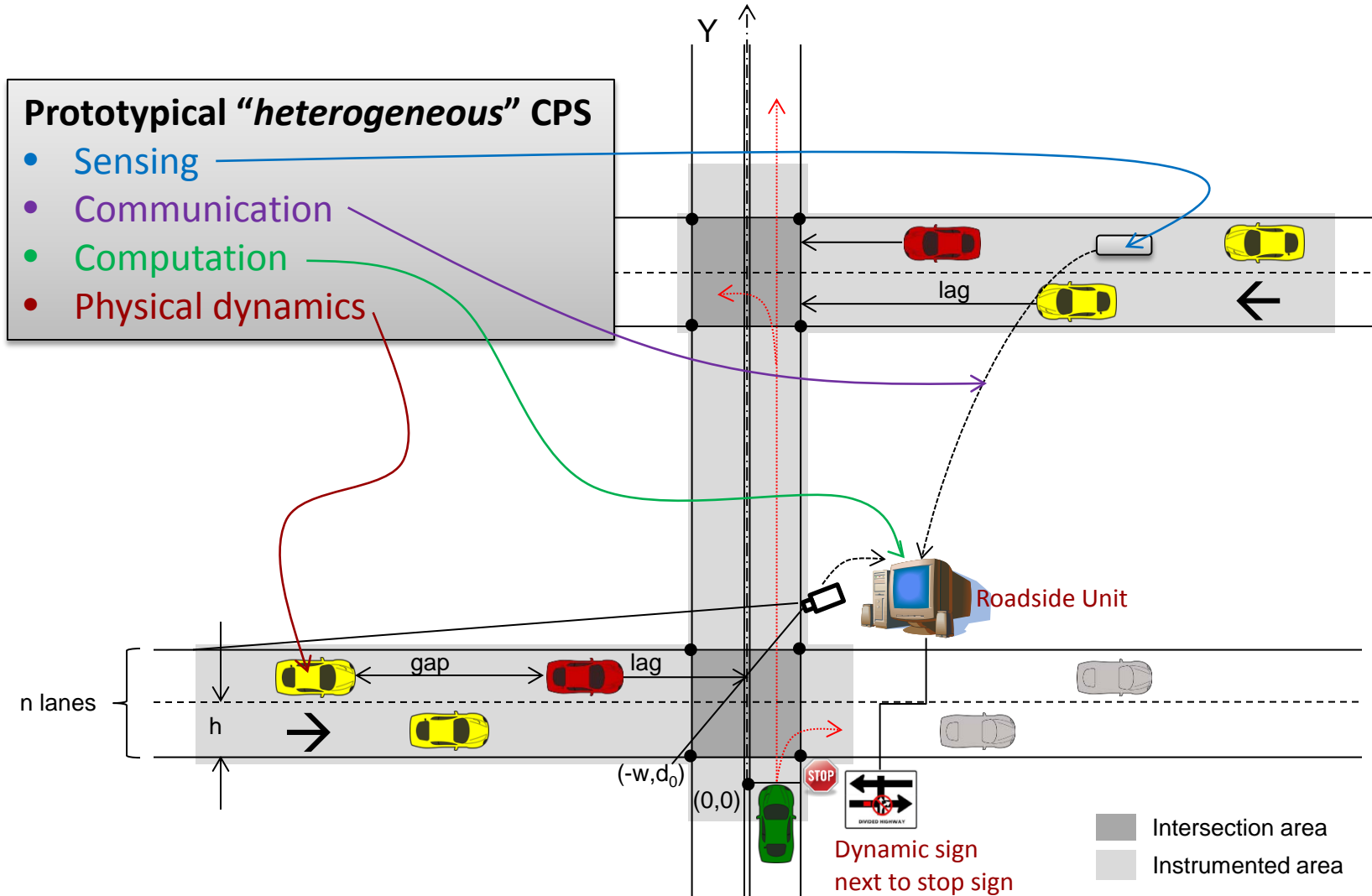
Joint work with

[Ajinkya Bhave](#), [Bruce Krogh](#), [David Garlan](#), [Sarah Loos](#), [Andre Platzer](#), [Ivan Ruchkin](#), [Bradley Schmerl](#)

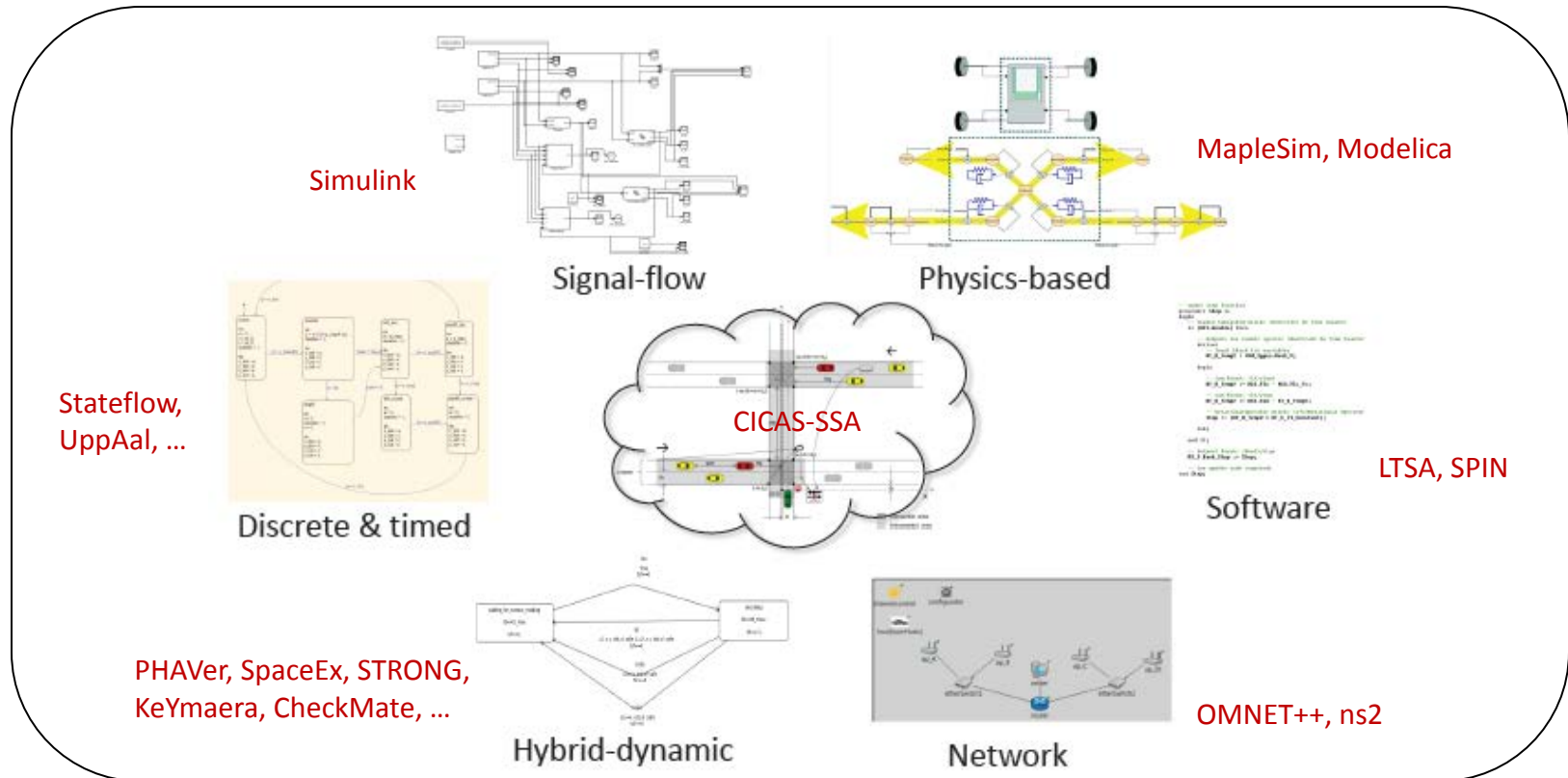
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CPS are *heterogeneous*

Example: CICAS-SSA*



Heterogeneity in Models & Analysis Tools



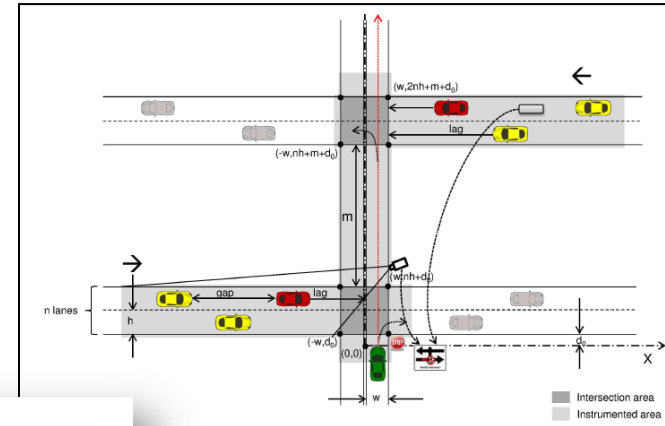
Challenges:

- *No “universal” modeling formalism* that can capture everything.
- Each model represents *some* design aspect well, but not the others.
- Models make (*interdependent*) *simplifying assumptions*.
- Different *tools* leverage different properties, *work only with their formalism*.

How do we ensure correctness of the system without a unifying formalism?

Architectural Modeling of CICAS

CICAS: Architectural modeling depicting interacting run-time components and connectors

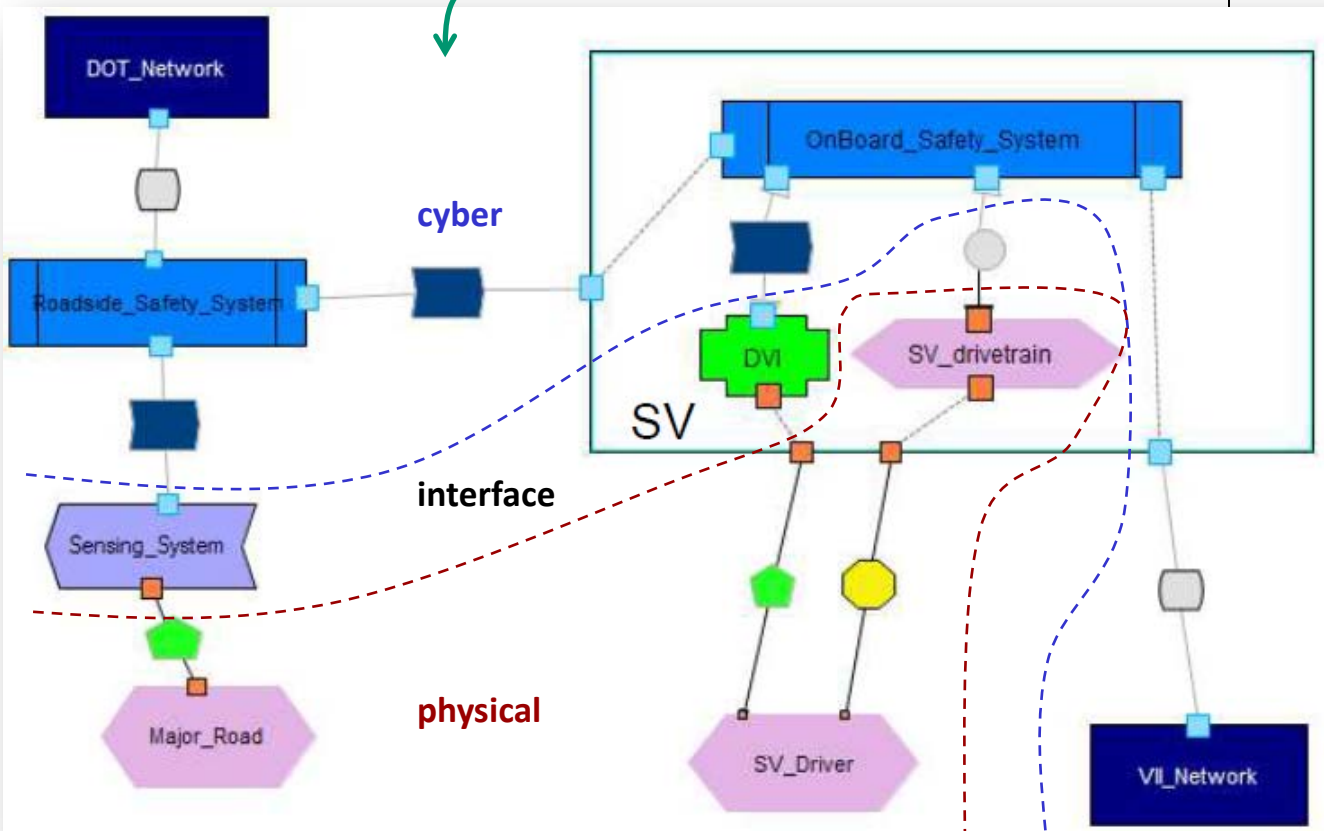


CICAS: Actual System



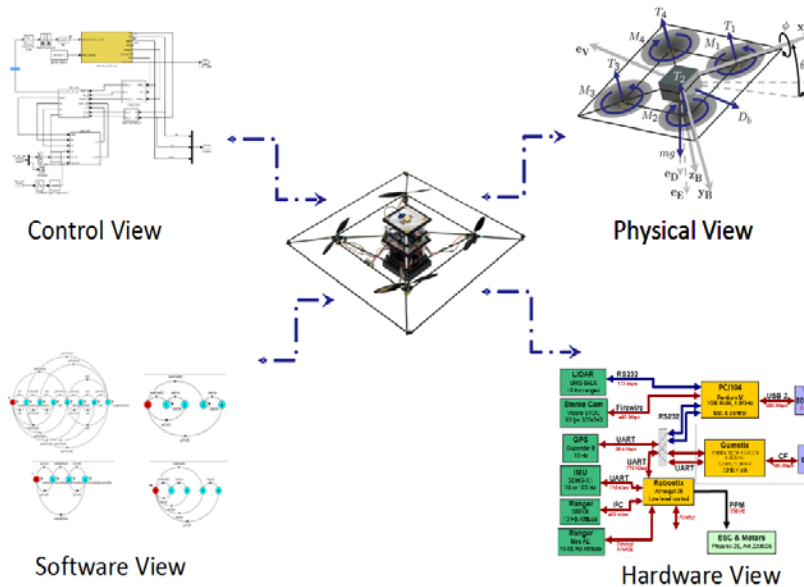
– *Visual representation with unambiguous (but very basic) semantics*

(no behavior info.)

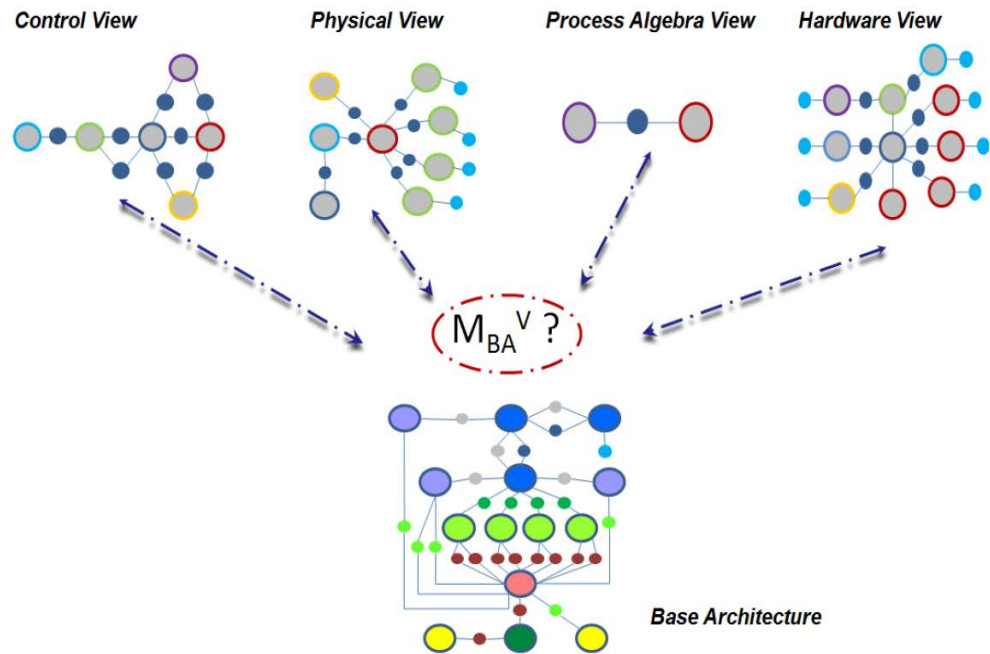


★ Heterogeneous models as arch. views

Example: STARMAC quadrotor



Models as architectural views



Structural consistency using graph morphisms

Augmenting Software Architectures with Physical Components

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ERTS² '10

View Consistency in Architectures for Cyber-Physical Systems

ICCPs '11

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★ Ensures consistent functional deployment in model subcomponents

★ Heterogeneous abstraction/implication

1. Define *behavior relations* R_i between abstract domains B_i and detailed domain B_0
2. Extend notions of abstraction/implication to heterogeneous domains via these R_i

Heterogeneous Abstraction

$M_0 \sqsubseteq^{R_1} M_1$, if (behavior set overapproximation via R_1)

$$\llbracket M_0 \rrbracket^{B_0} \subseteq R_1^{-1}(\llbracket M_1 \rrbracket^{B_1}). \quad \text{A}$$

Heterogeneous Specification Implication

$S_1 \Rightarrow^{R_1} S_0$, if (behavior set underapproximation via R_1)

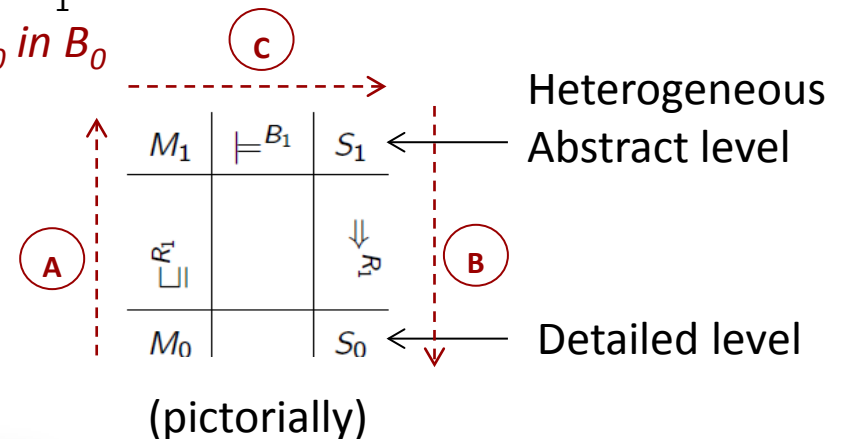
$$R_1^{-1}(\llbracket S_1 \rrbracket^{B_1}) \subseteq \llbracket S_0 \rrbracket^{B_0}. \quad \text{B}$$

3. If M_1 abstracts M_0 and S_1 stronger than S_0 via R_1 then M_1 satisfies S_1 in B_1 implies M_0 satisfies S_0 in B_0

Heterogeneous Verification

If $M_0 \sqsubseteq^{R_1} M_1$, $M_1 \models^{B_1} S_1$ and $S_1 \Rightarrow^{R_1} S_0$, then $M_0 \models^{B_0} S_0$. C

(in words)



Heterogeneous Verification of Cyber-Physical Systems
using Behavior Relations

HSCC '12

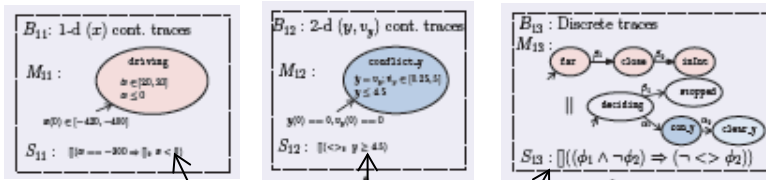
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★ Enables heterogeneous verification

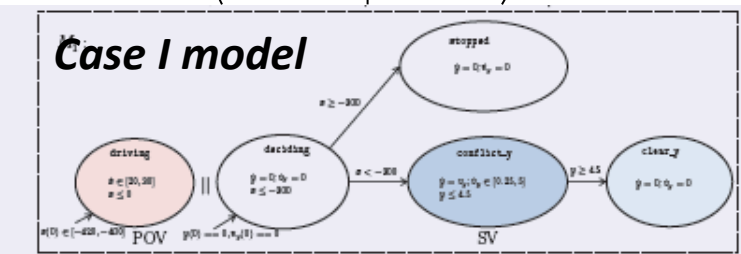
★ Hierarchical heterogeneous verification of CICAS-SSA

of CICAS-SSA



Conjunctive breakdowns

- M_{1i} individually cover M_1
- S_{1i} together imply S_1

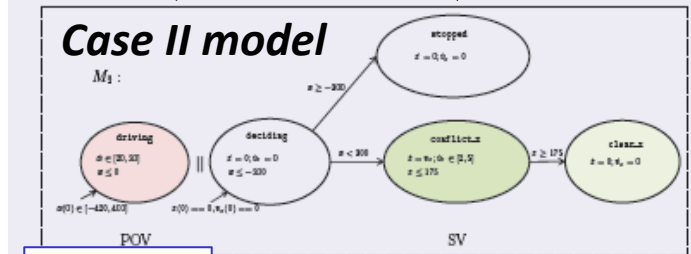


No right turn

B_1 : 3-d $\{(x, y, v_y)\}$ hybrid traces
 S_1 : $\square \neg (x == 0 \wedge 0 < y < 4.5)$

R_1 : equal values along x, y, v_y ; z, v_z zero over all time

$R_1^{-1}([\![S_1]\!]^{B_1})$: $\square ((\neg(x == 0 \wedge 0 < y < 4.5)) \wedge (z == 0 \wedge v_z == 0))$



No straight crossing

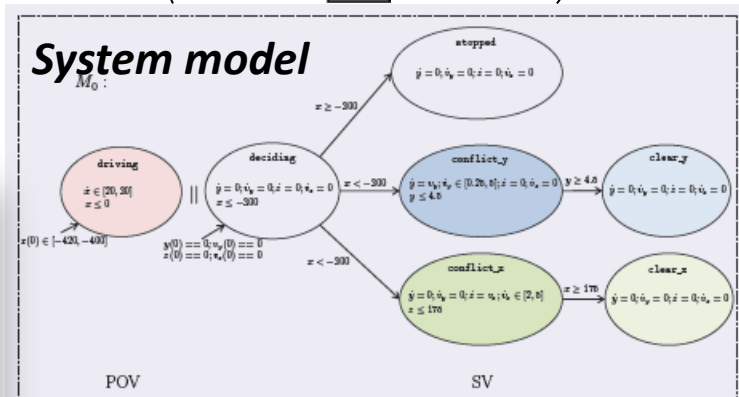
B_2 : 3-d $\{(x, z, v_z)\}$ hybrid traces
 S_2 : $\square \neg (x == 0 \wedge 0 < z < 170)$

R_2 : equal values along x, z, v_z ; y, v_y zero over all time

$R_2^{-1}([\![S_2]\!]^{B_2})$: $\square ((\neg(x == 0 \wedge 0 < z < 170)) \wedge (y == 0 \wedge v_y == 0))$

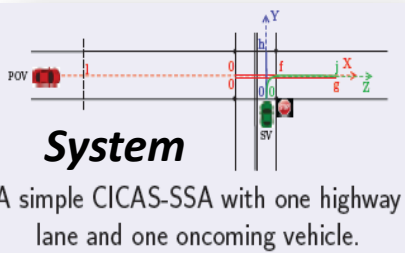
Disjunctive breakdown

- M_i together cover M_0
- S_i individually imply S_0

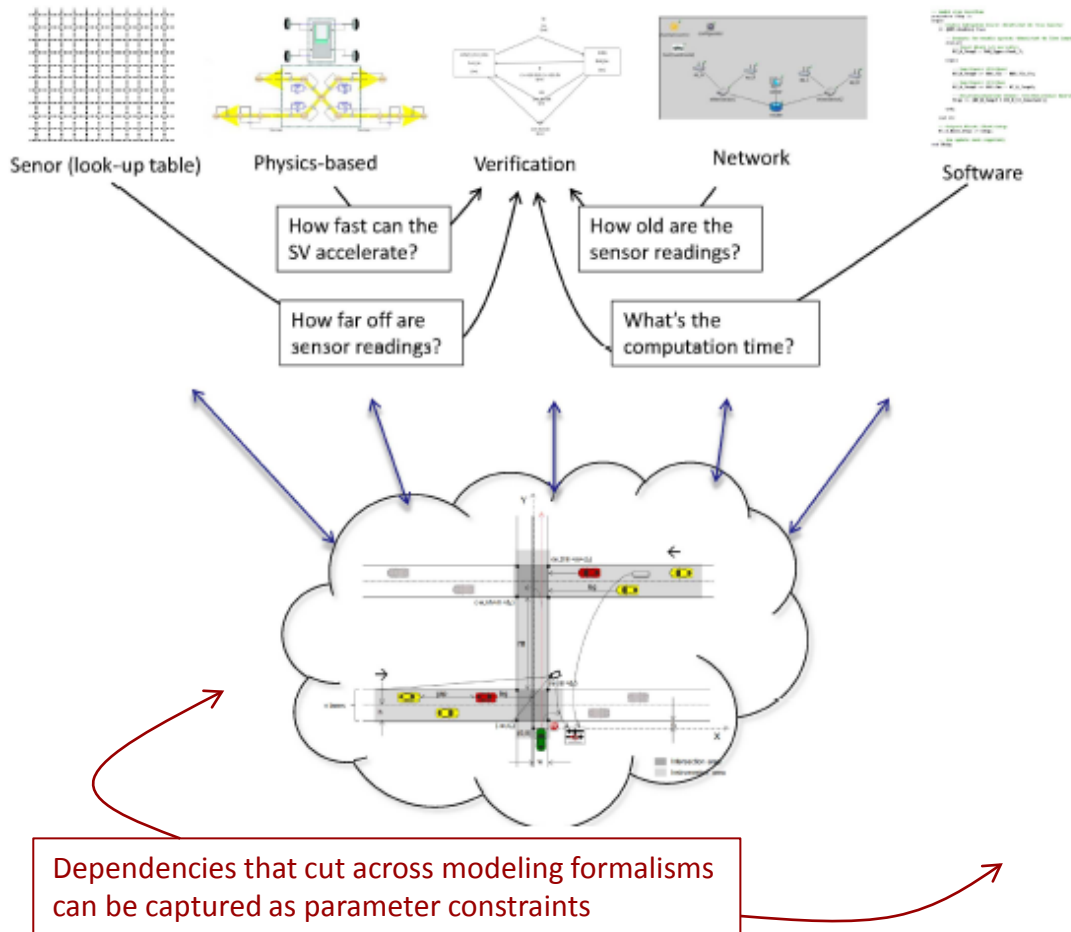


B_0 : 5-d $\{(x, y, v_y, z, v_z)\}$ hybrid traces
 S_0 : $\square \neg ((x == 0 \wedge 0 < y < 4.5) \vee (x == 0 \wedge 0 < z < 170))$

★ 5-d hybrid verification simplified to 1-d, 2-d continuous and discrete verifications



★ Semantic assumptions as parameter constraints



Problem

- Semantic interdependencies cut across modeling formalisms.
- Consistency needs to be ensured to guarantee system verification.

Challenge

- Interdependencies need to be formally represented
- Representation needs to be universal to all modeling formalisms

Approach

- Identify model and spec. parameters explicitly
- Model interdependencies as an *auxiliary constraint*
- Find *effective constraint* on given model/spec. parameters using projection (existential quantification)
- Prove consistency in SMT solvers or theorem provers

Using Parameters in Architectural Views to Support Heterogeneous Design and Verification

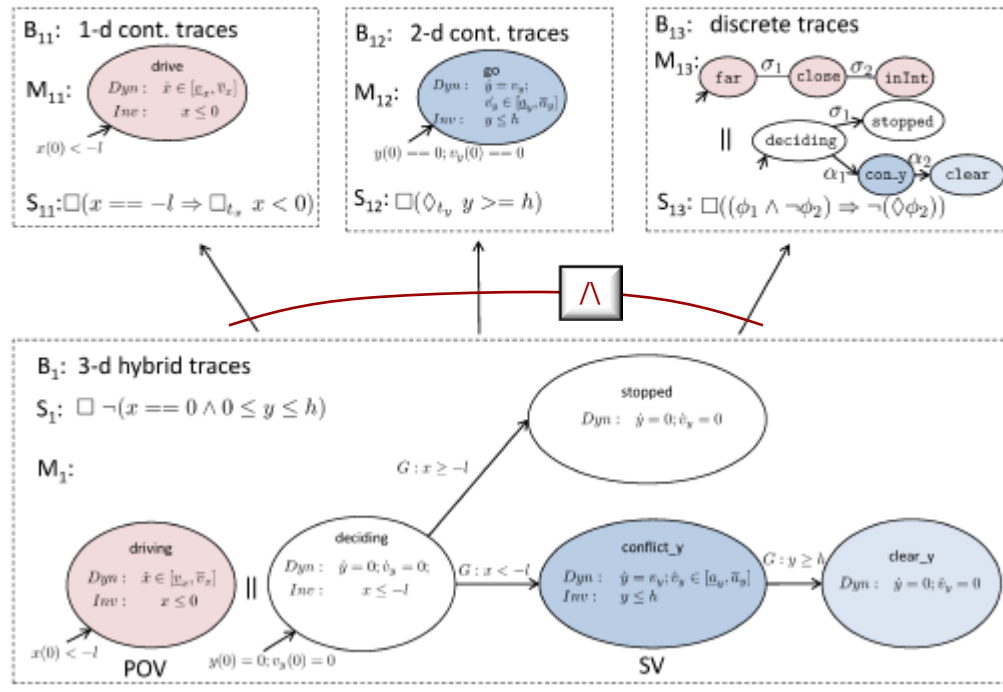
CDC '11

Akshay Rajhans[†], Ajinkya Bhav[†], Sarah Loos[‡], Bruce H. Krogh[†], André Platzer[‡], David Garlan[‡]

★ Ensures semantic (parameter) consistency using external SMT solvers or provers

★ Parametric verification of CICAS

Parameterized models and specifications



1. Explicitly identify model parameters e.g. *speed limits, intersection geometry, minimum acceleration*, and spec. parameters, e.g., *POV min. time-to-intersection, SV max. time-to-clear-intersection*

2. Model interdependencies as an auxiliary constraint e.g., those dictated by *speed limits, newton's laws* and *intersection geometry* on *time-to-intersection*, ...

3. Project global constraints and interdependencies (aux. constraint) onto local sets of parameters

Heterogeneous Verification of Cyber-Physical Systems using Behavior Relations

HSCC '12

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★ Proved semantic consistency in theorem prover KeYmaera