

Model-Based Design of Connected and Autonomous Vehicles

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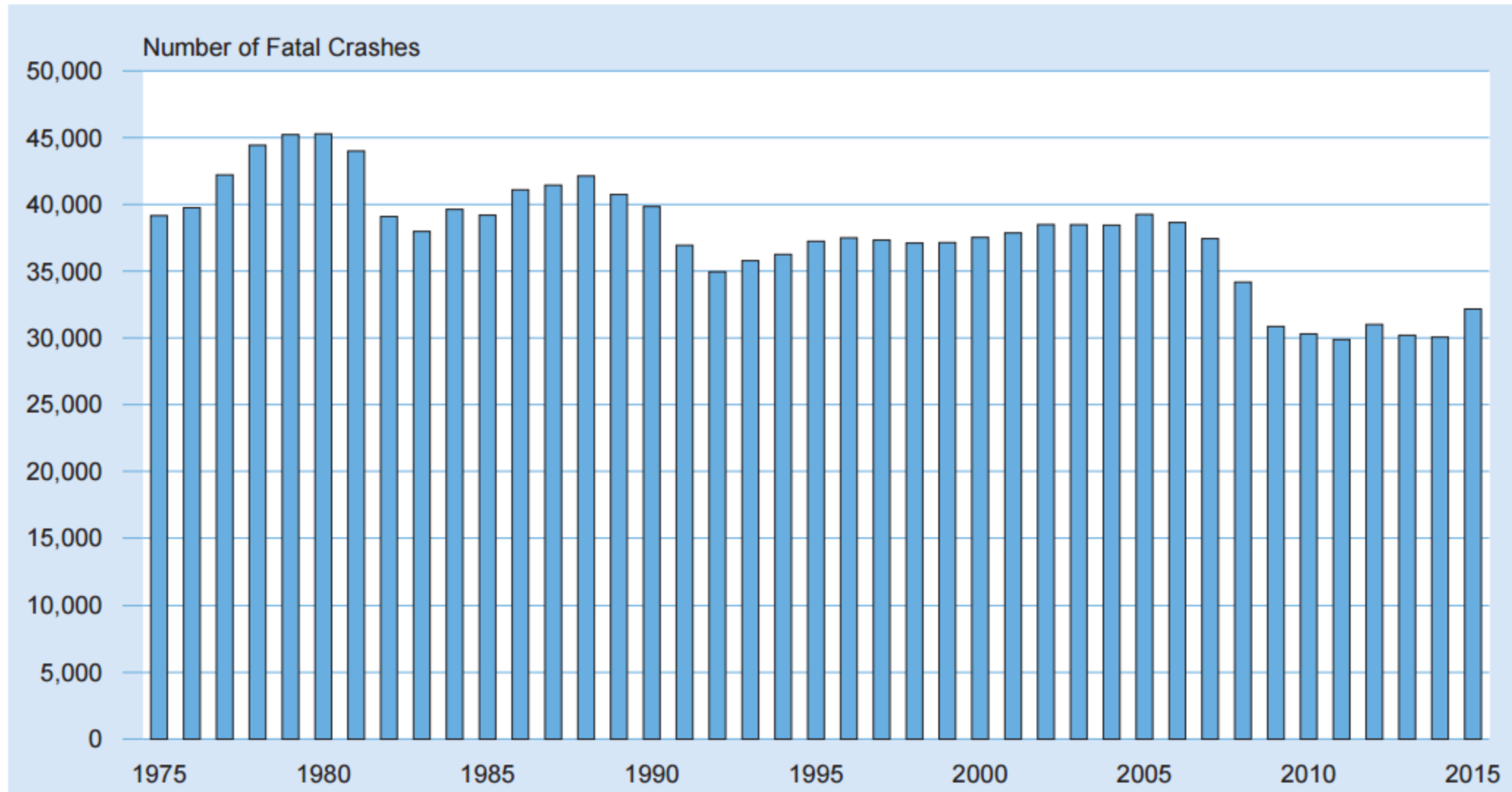
2nd IEEE Summer School on Connected and Autonomous Vehicles
May 19, 2017

How do we design safe and reliable
cyber-physical systems ?

Model-based design (MBD)

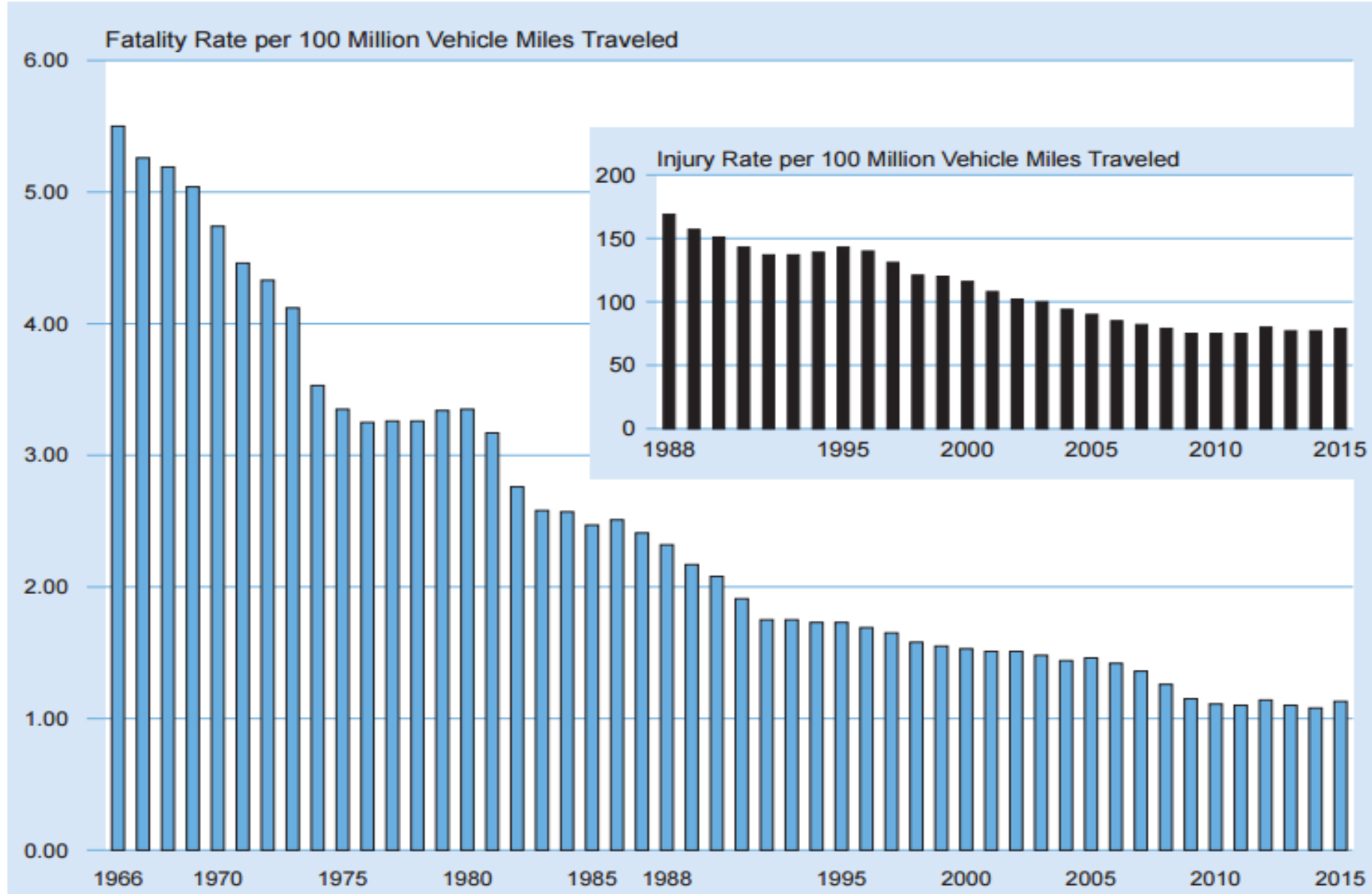
- Analyze and understand the requirements specification
- Develop computational model(s) of the system
 - Check the model against the real system
 - “are you are building the right thing?” (validation)
 - Check the model against specifications
 - “are you building it right?” (verification)
- Build a prototype
 - test the prototype in the actual working environment
- Production

Fatal Crashes, 1975-2015



<https://crashstats.nhtsa.dot.gov/Api/Public/ViewPublication/812384>

Motor Vehicle Fatality and Injury Rates per 100 Million Vehicle Miles Traveled, 1966-2015



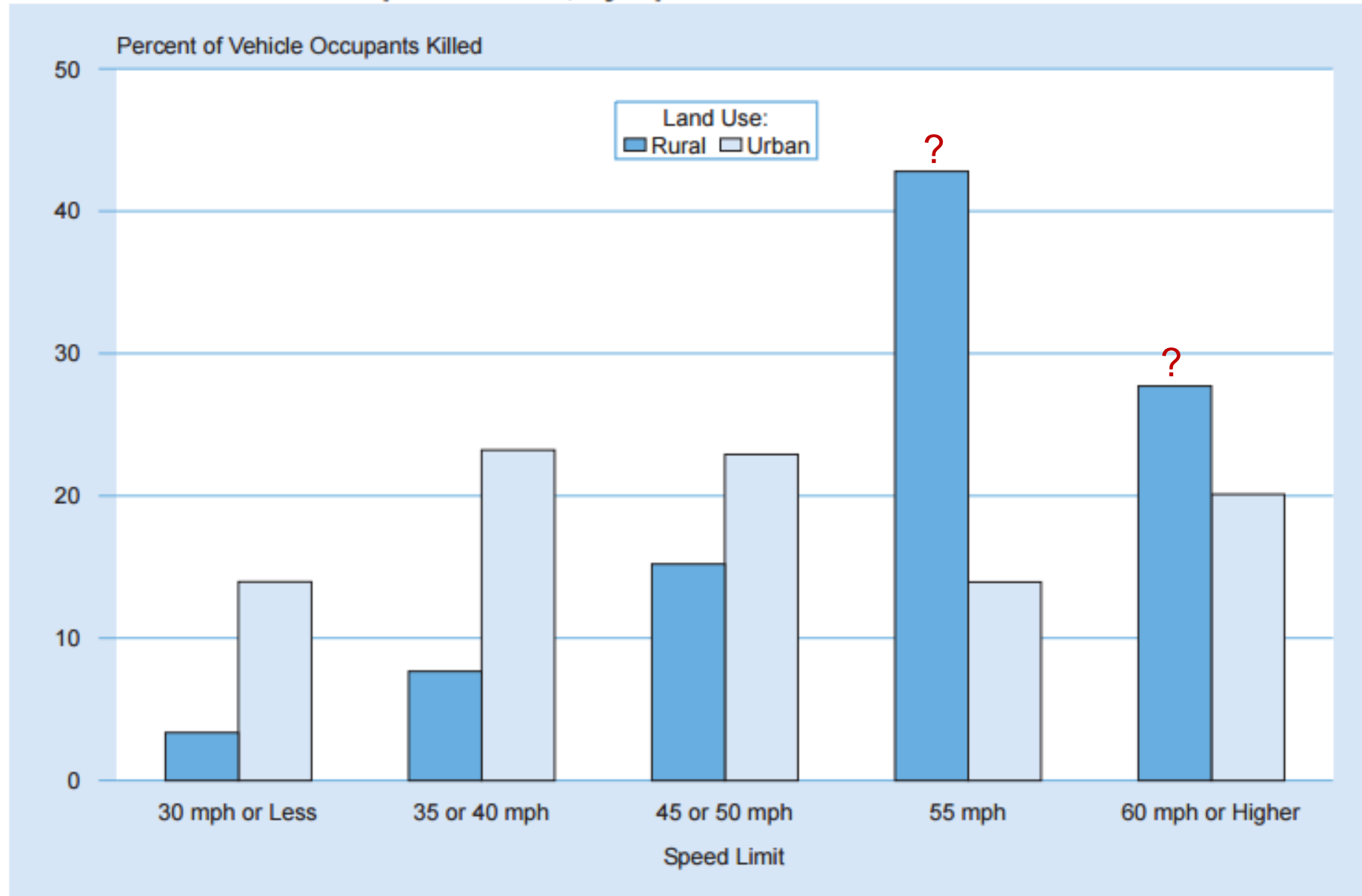
<https://crashstats.nhtsa.dot.gov/Api/Public/ViewPublication/812384>

Vehicles Involved in Fatal Crashes by Speed Limit and Land Use

Speed Limit	Land Use						Total	
	Rural		Urban		Unknown			
	Number	Percent	Number	Percent	Number	Percent	Number	Percent
30 mph or less	707	15.8	3,033	67.9	725	16.2	4,465	100.0
35 or 40 mph	1,707	20.6	5,523	66.5	1,071	12.9	8,301	100.0
45 or 50 mph	3,506	35.9	5,374	55.0	890	9.1	9,770	100.0
55 mph	9,743	74.8	2,928	22.5	351	2.7	13,022	100.0
60 mph or higher	6,600	60.0	4,152	37.7	254	2.3	11,006	100.0
No Statutory Limit	113	33.6	177	52.7	46	13.7	336	100.0
Unknown	629	31.1	1,187	58.7	207	10.2	2,023	100.0
Total	23,005	47.0	22,374	45.7	3,544	7.2	48,923	100.0

<https://crashstats.nhtsa.dot.gov/Api/Public/ViewPublication/812384>

Percent of Vehicle Occupants Killed, by Speed Limit and Land Use



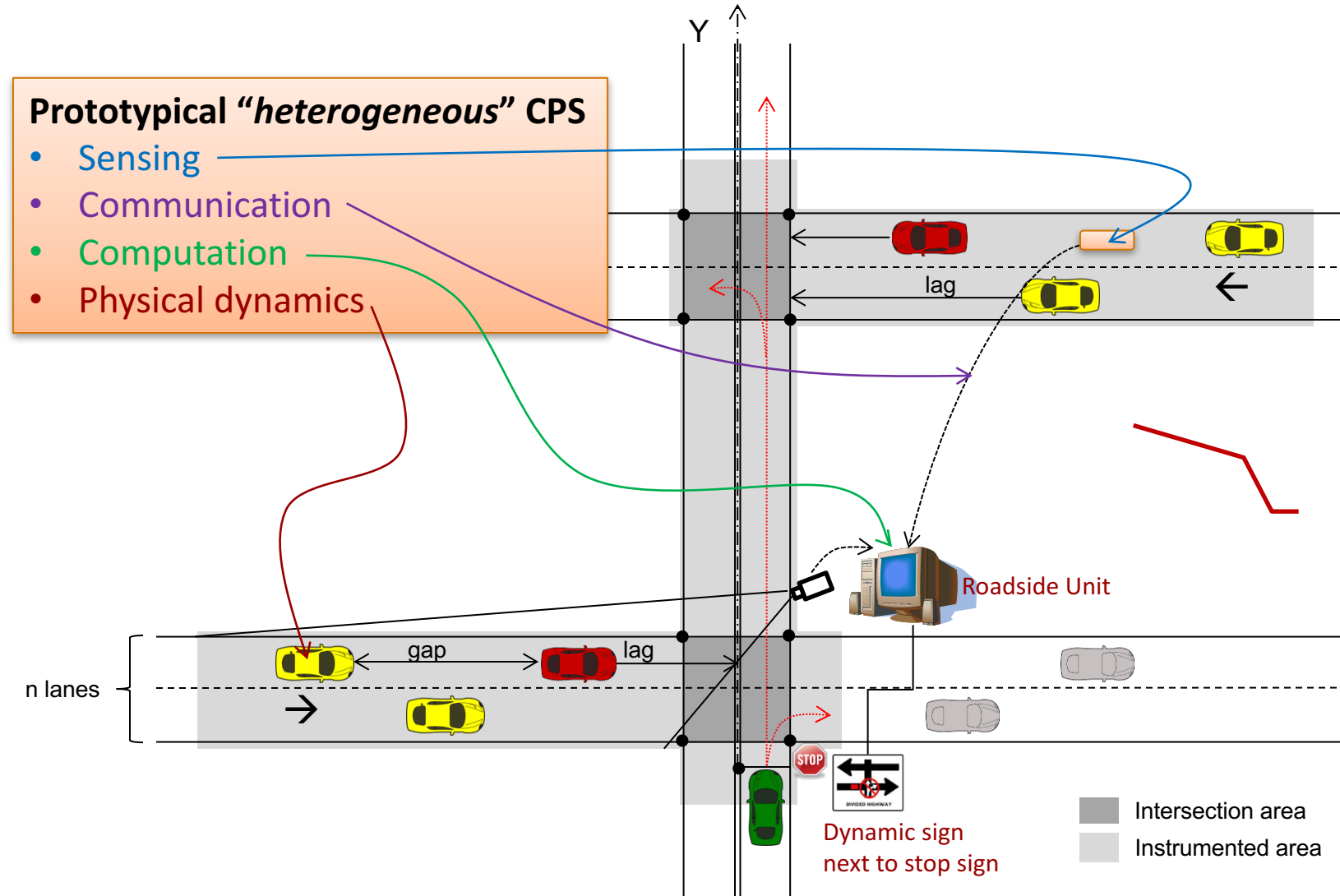
<https://crashstats.nhtsa.dot.gov/Api/Public/ViewPublication/812384>

Cooperative Intersection Collision Avoidance System: Stop-Sign Assist (CICAS-SSA)

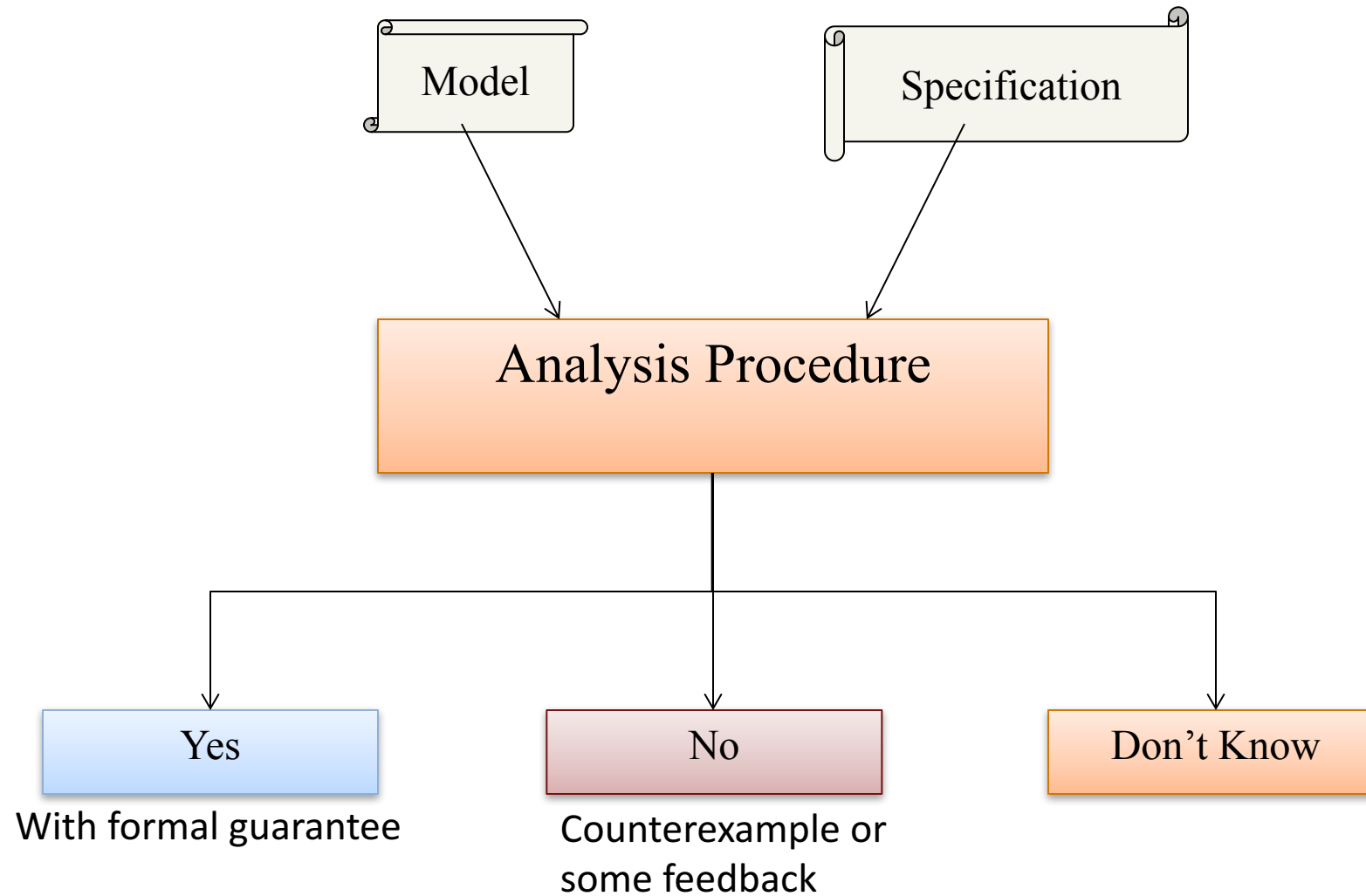


Figure 1: Plan view of a typical instrumented rural four lane expressway intersection. Sensors are radar (yellow triangles indicate field of view and) scanning lidar (orange semicircles); all data is sent from sensor processors to the main central processor.

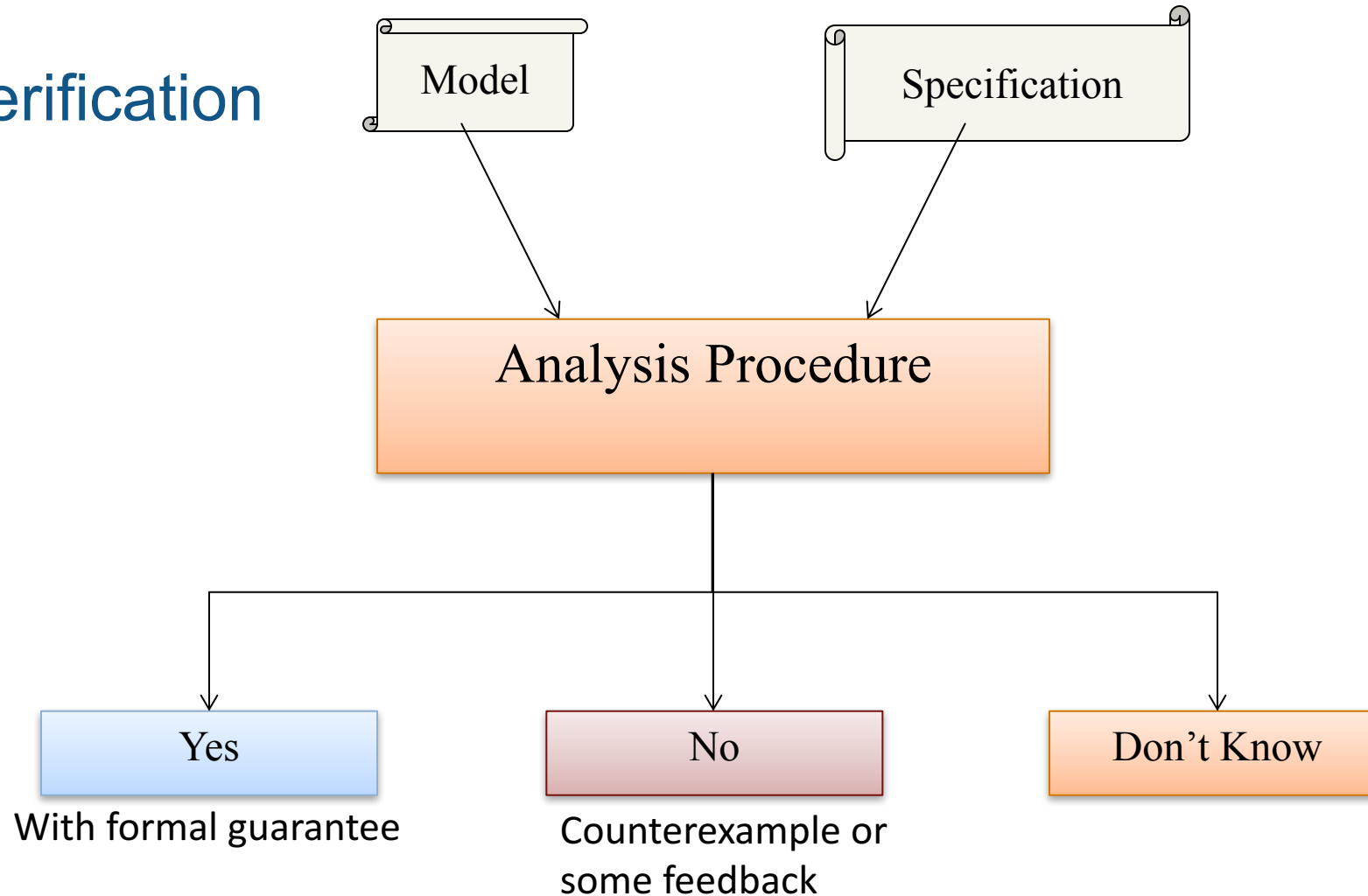
CICAS-SSA Schematic



Formal Verification

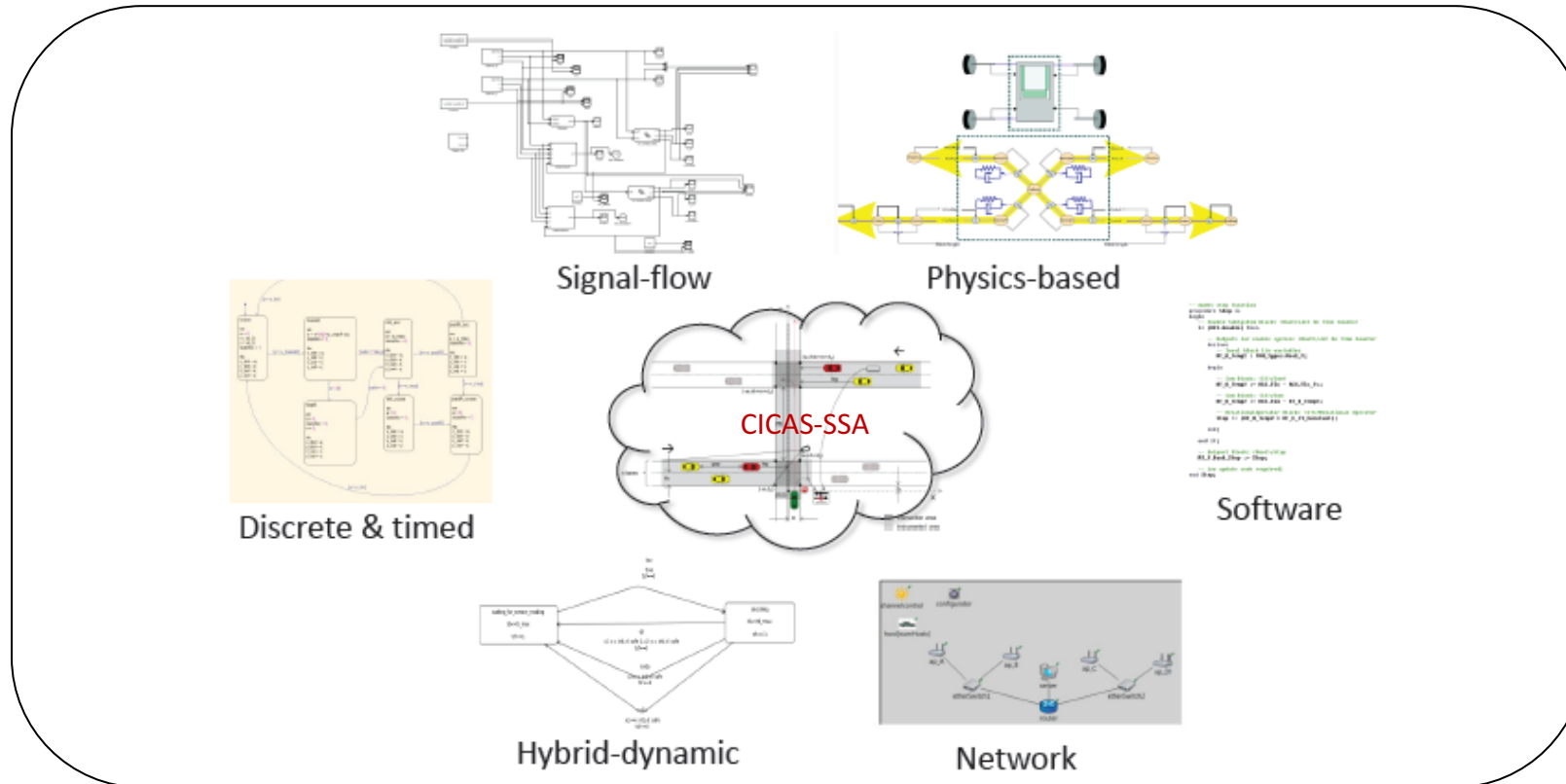


Formal Verification



There is no system model!
But there are models...

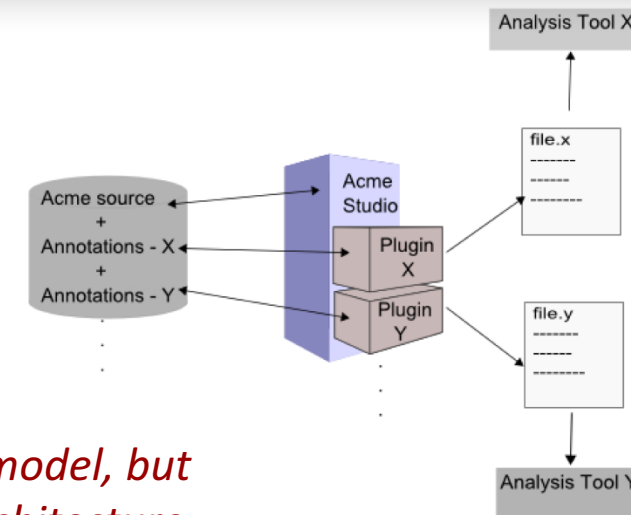
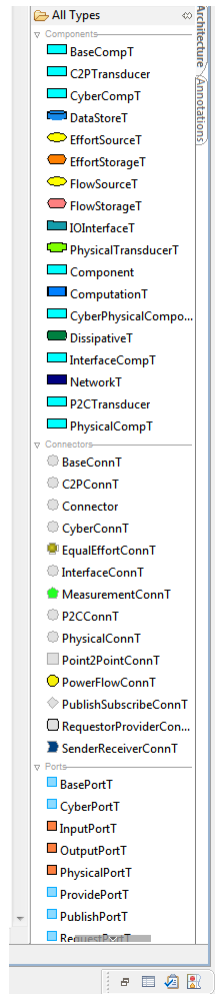
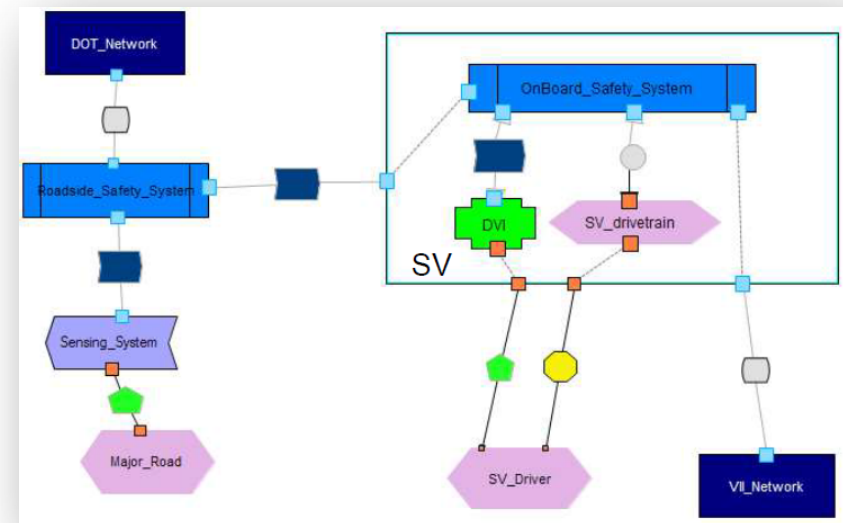
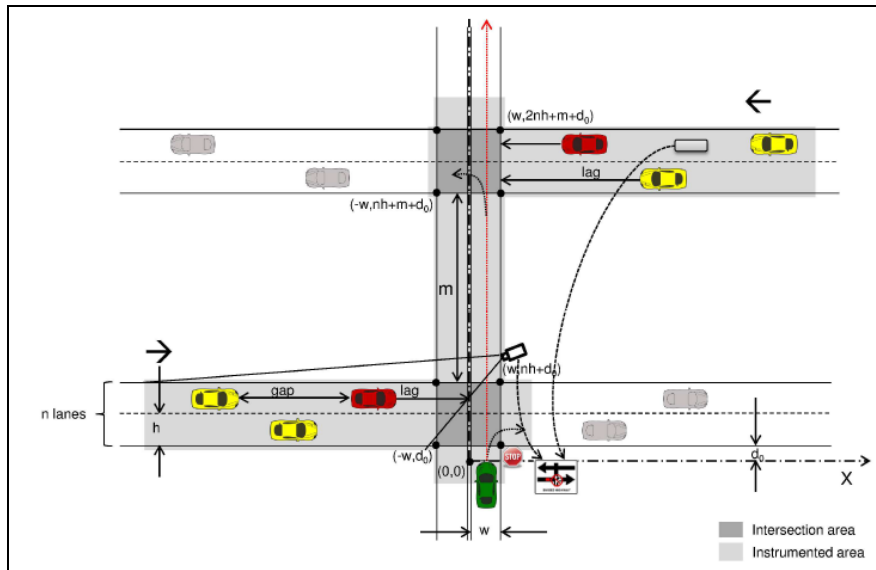
Heterogeneity in modeling formalisms and analysis techniques



- *Different formalisms suited for different aspects of system design*
- Each model represents *some* design aspect well
- Models make *interdependent assumptions*
- *Tools work only with their formalisms*

How do we ensure correctness of the system?

Cyber-Physical System Architecture



MPM '09

ECEASST

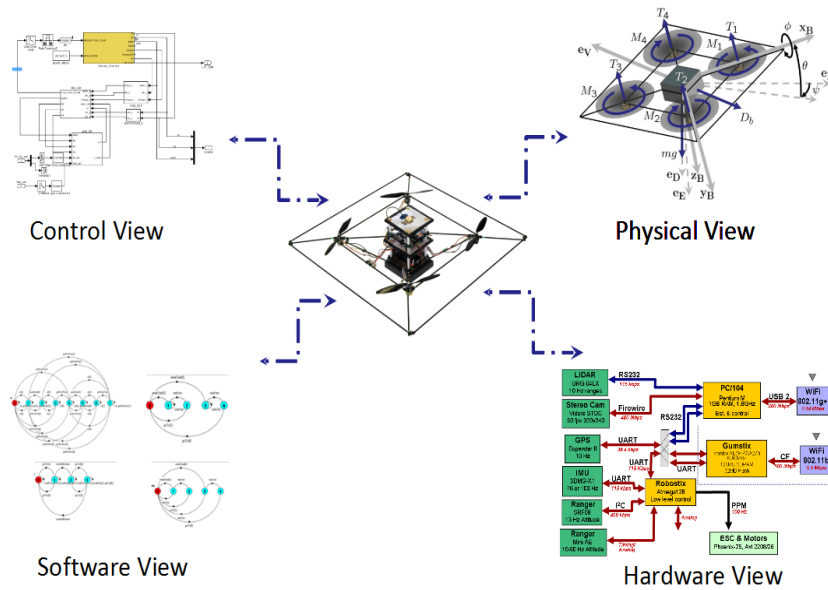
An Architectural Approach to the Design and Analysis of Cyber-Physical Systems
 Akshay Rajhans¹, Shang-Wen Cheng², Bradley Schmerl², David Garlan², Bruce H. Krogh¹, Clarence Agbi¹ and Ajinkya Bhawe¹



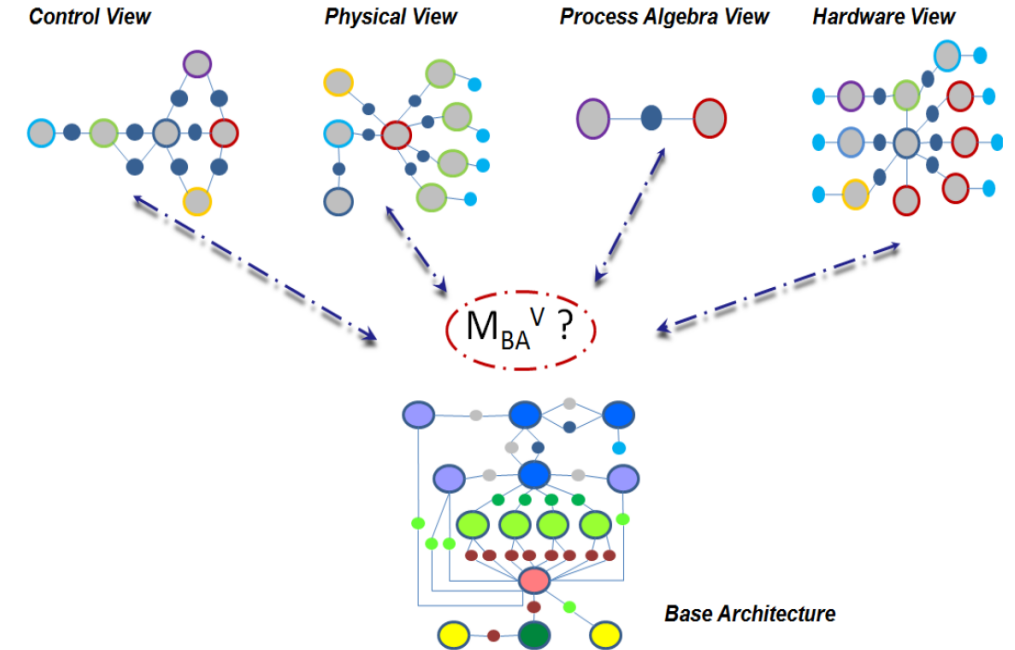
There is no system model, but there is a system architecture

CPS architectural style palette in AcmeStudio

Architectural views



Models as architectural views



Structural consistency using graph morphisms

Augmenting Software Architectures with Physical Components

Ajinkya Bhavé¹, David Garlan², Bruce H. Krogh¹, Akshay Rajhans¹, Bradley Schmerl²

ERTS² '10

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View Consistency in Architectures for Cyber-Physical Systems

ICCPs '11

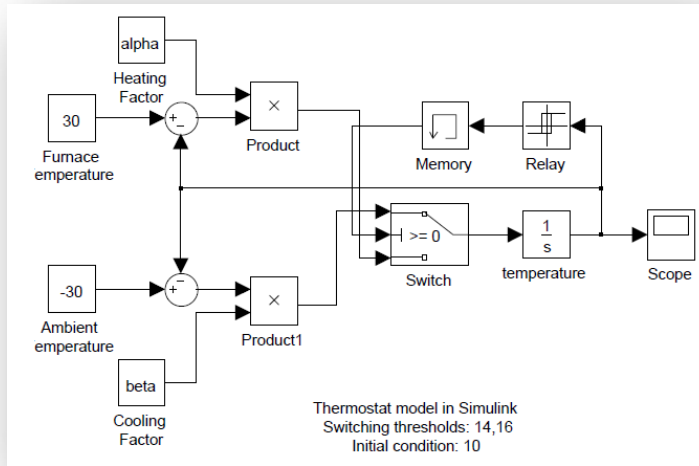
Ajinkya Bhavé, Bruce H. Krogh

David Garlan, Bradley Schmerl

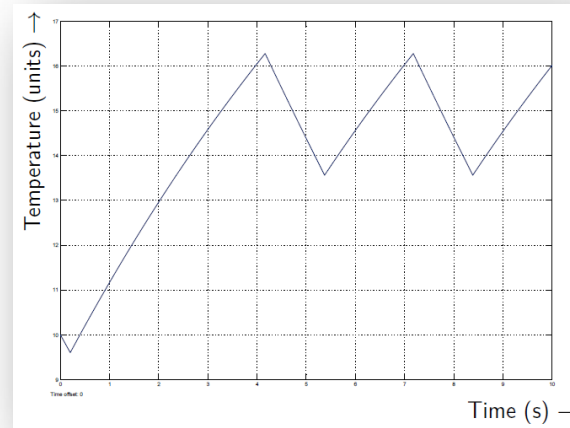


"Model structure vs system structure"
Analysis: Consistency, completeness

Semantic domains of models and specifications



Model M

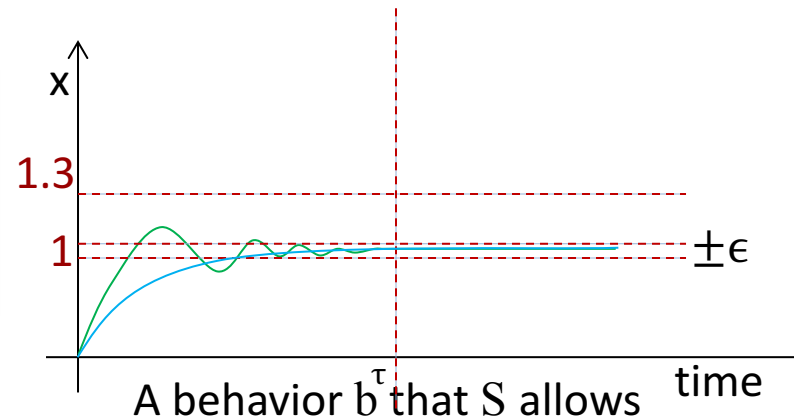


A behavior b that M exhibits

1) “overshoot is no more than 1.3 units and settling time is less than τ ”

2) $\Box(x < 1.3) \wedge \Diamond_{\tau}(x \in [1 \pm \epsilon])$

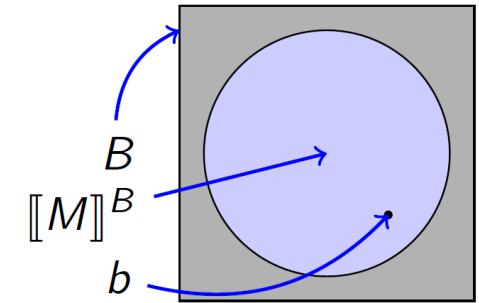
Specification S



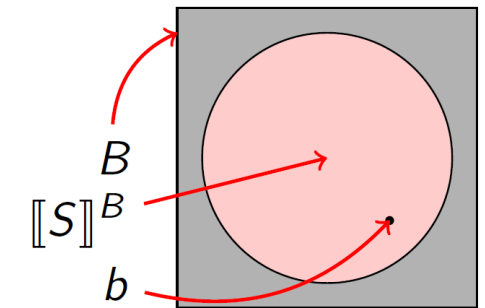
A behavior b that S allows

$$M \models^B S$$

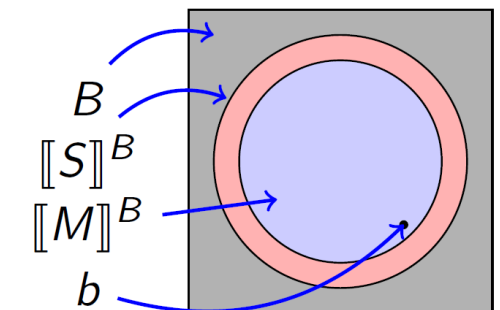
$$[M]^B \subseteq [S]^B$$



$[M]^B$: “semantic interpretation” of M in a behavior domain B

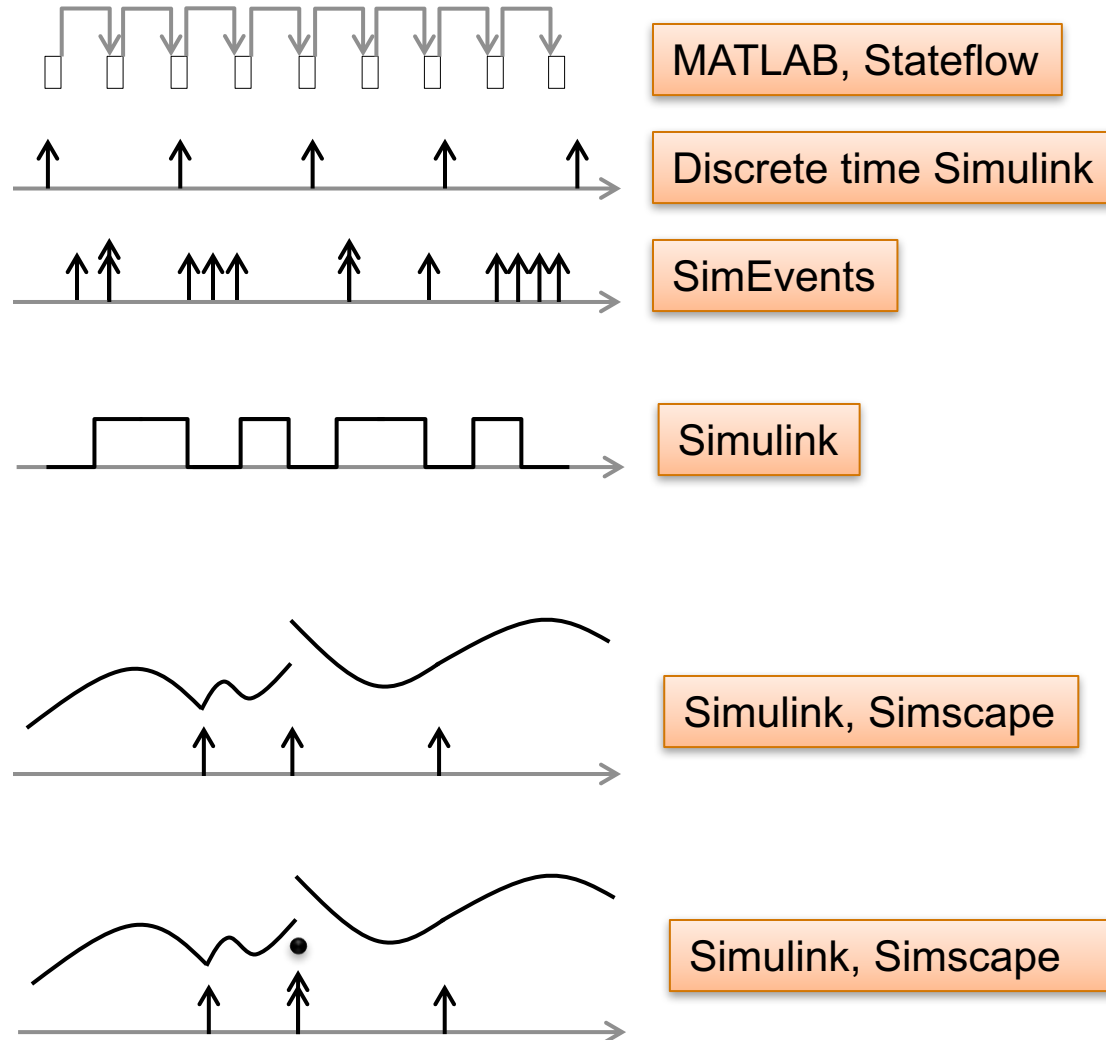


$[S]^B$: “semantic interpretation” of S in B



The semantic domain of a dynamic system

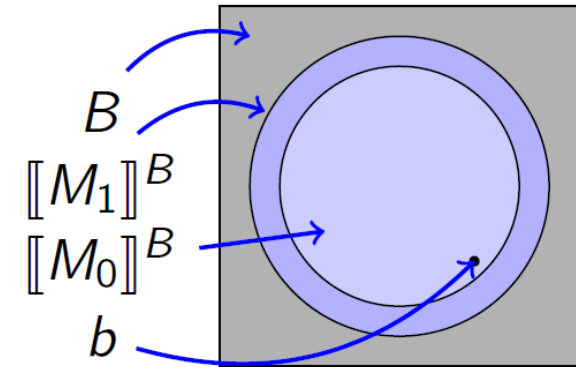
- Points, []
 - On \mathbf{N}
 - On $\mathbf{R} \times \mathbf{N}$
- Intervals, [> (< >, <]
 - On \mathbf{R}
- Hybrid point/interval
 - On \mathbf{R}
 - On $\mathbf{R} \times \mathbf{N}$



Abstraction and Implication

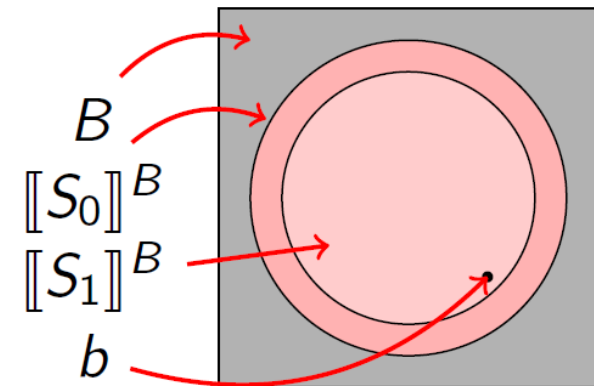
- Model M_1 abstracts M_0 in B , written $M_0 \sqsubseteq^B M_1$

$$\text{if } \llbracket M_0 \rrbracket^B \subseteq \llbracket M_1 \rrbracket^B$$



- Specification S_1 implies S_0 in B , written $S_1 \Rightarrow^B S_0$

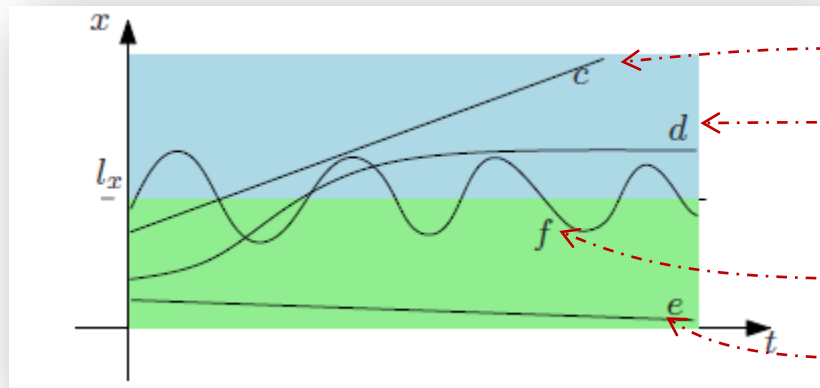
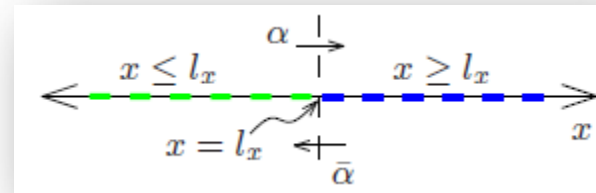
$$\text{if } \llbracket S_1 \rrbracket^B \subseteq \llbracket S_0 \rrbracket^B$$



Mappings between semantic domains via behavior relations

- *Approach*: Create “behavior relations” between domains

Example



$$R_1 \subseteq B_0 \times B_1$$

α

$\alpha \bar{\alpha} \alpha \bar{\alpha} \alpha \bar{\alpha} \dots$

ε

Given $R_1 \subseteq B_0 \times B_1$
set-based inverse map
 $R_1^{-1}(' \alpha ') = \{c, d, \dots\}$

B_0 : 1-d continuous trajectories in x

$$B_1 = \{\alpha, \bar{\alpha}\}^* \cup \{\alpha, \bar{\alpha}\}^\omega$$

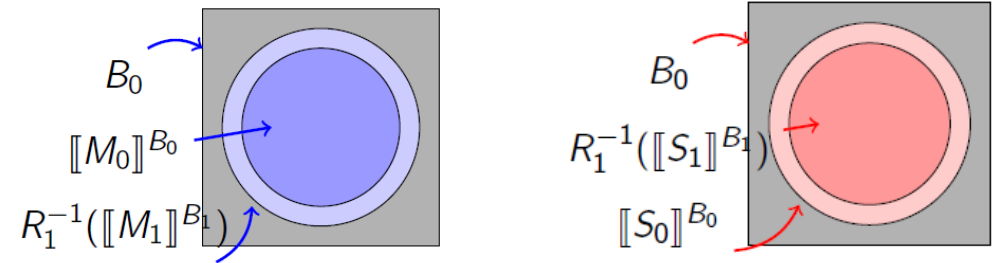
Heterogeneous Abstraction and Implication

■ Heterogeneous extensions of behavior-set inclusions

Heterogeneous Abstraction

$M_0 \sqsubseteq^{R_1} M_1$, if

A $\llbracket M_0 \rrbracket^{B_0} \subseteq R_1^{-1}(\llbracket M_1 \rrbracket^{B_1}).$



Heterogeneous Specification Implication

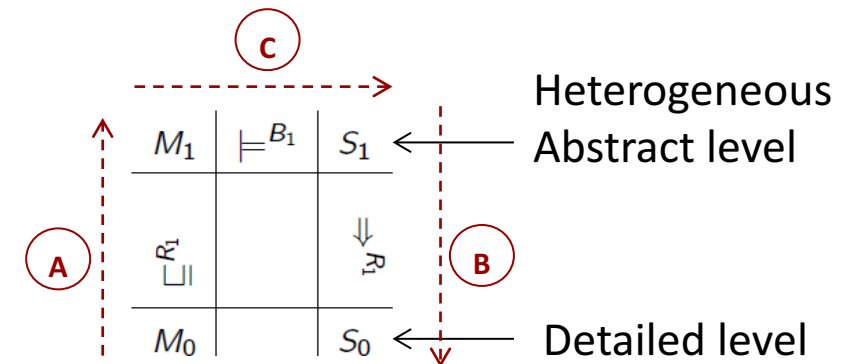
$S_1 \Rightarrow^{R_1} S_0$, if

B $R_1^{-1}(\llbracket S_1 \rrbracket^{B_1}) \subseteq \llbracket S_0 \rrbracket^{B_0}.$

Heterogeneous Verification

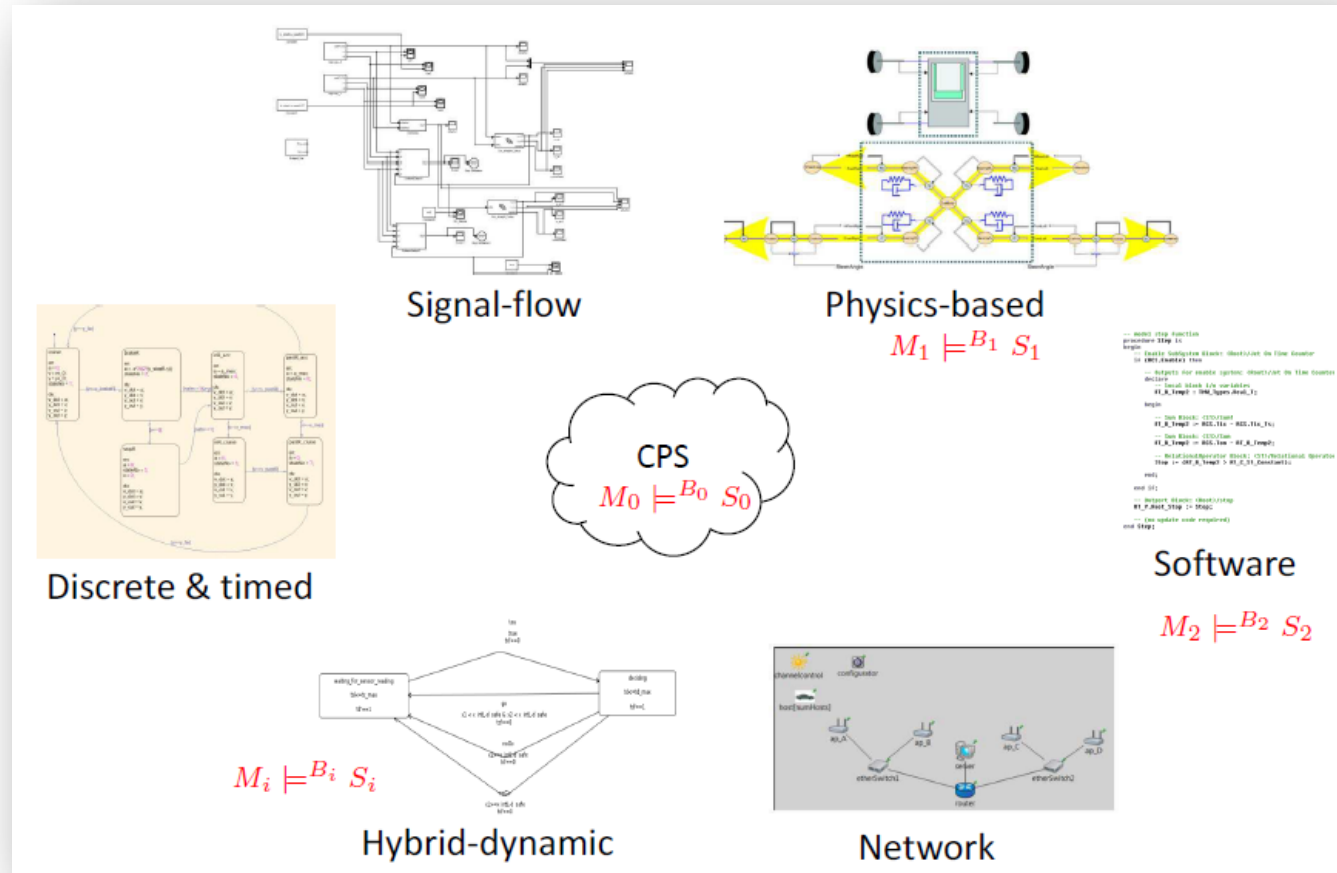
If $M_0 \sqsubseteq^{R_1} M_1$, $M_1 \models^{B_1} S_1$ and $S_1 \Rightarrow^{R_1} S_0$,
then $M_0 \models^{B_0} S_0$. **C**

(in words)



(pictorially)

Multi-model Verification Problem



How do we use $M_1 \models^{B_1} S_1, \dots, M_n \models^{B_n} S_n$ to infer $M_0 \models^{B_0} S_0$?

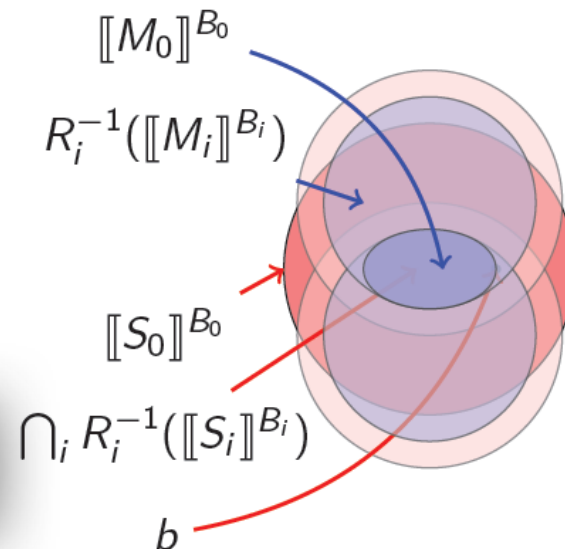
Multi-model conjunctive and disjunctive heterogeneous verification

Conjunctive specification implication

Given behavior relations $R_i \subseteq B_0 \times B_i$, a set of specifications S_1, \dots, S_n *conjunctively imply* S_0 if $\bigcap_i R_i^{-1}(\llbracket S_i \rrbracket^{B_i}) \subseteq \llbracket S_0 \rrbracket^{B_0}$.

Typical use case

- Each model captures a different aspect
- Specs pertain to only the relevant one

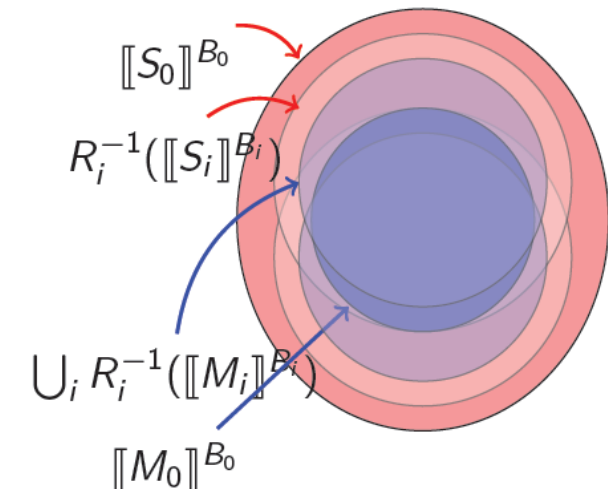


Model coverage (disjunctive abstraction)

Given behavior relations $R_i \subseteq B_0 \times B_i$, a set of models M_1, \dots, M_n *cover* M_0 if $\llbracket M_0 \rrbracket^{B_0} \subseteq \bigcup_i R_i^{-1}(\llbracket M_i \rrbracket^{B_i})$.

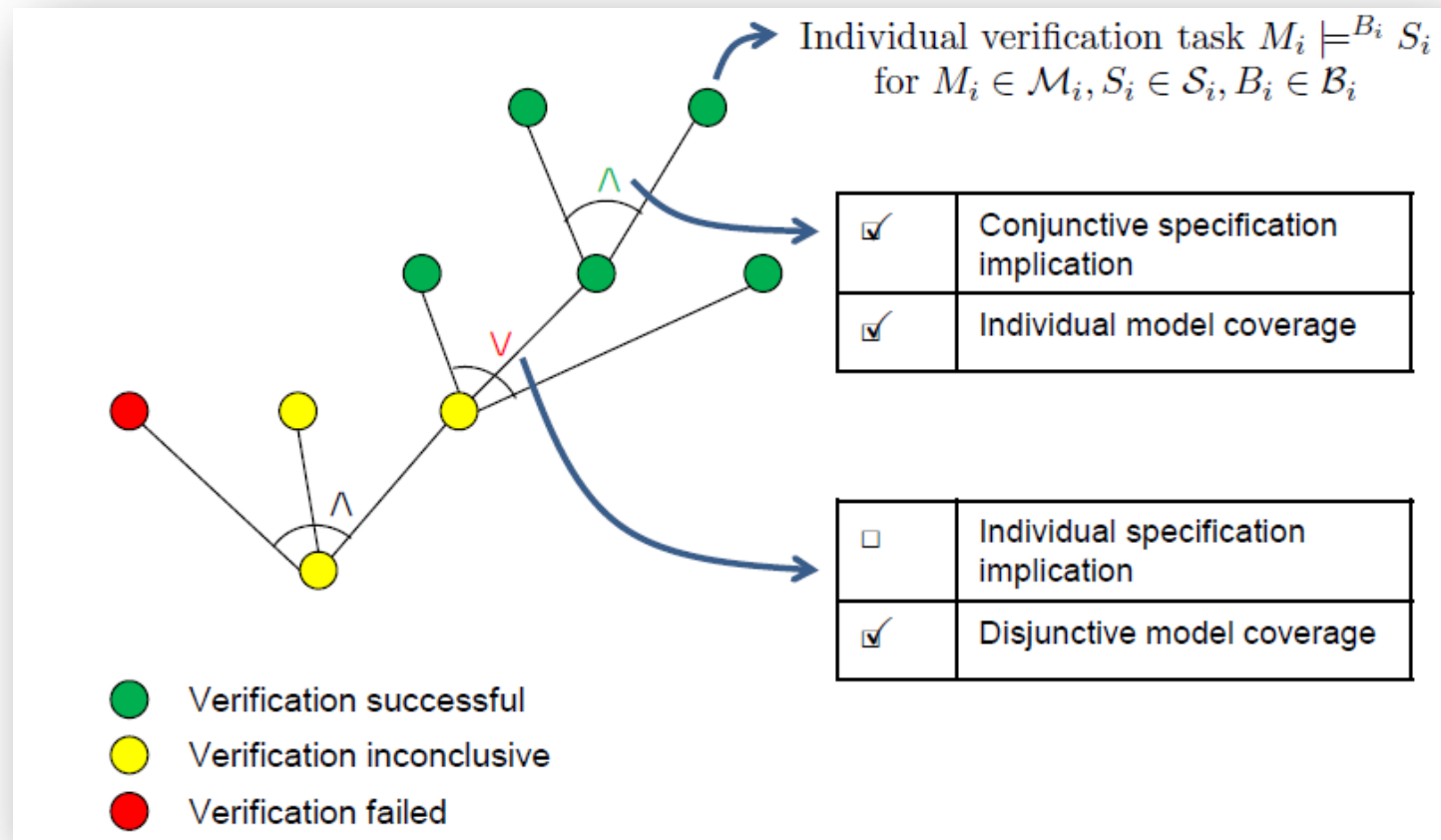
Typical use case

- Each model captures a different subset of behaviors, e.g., a specific nondeterministic choice



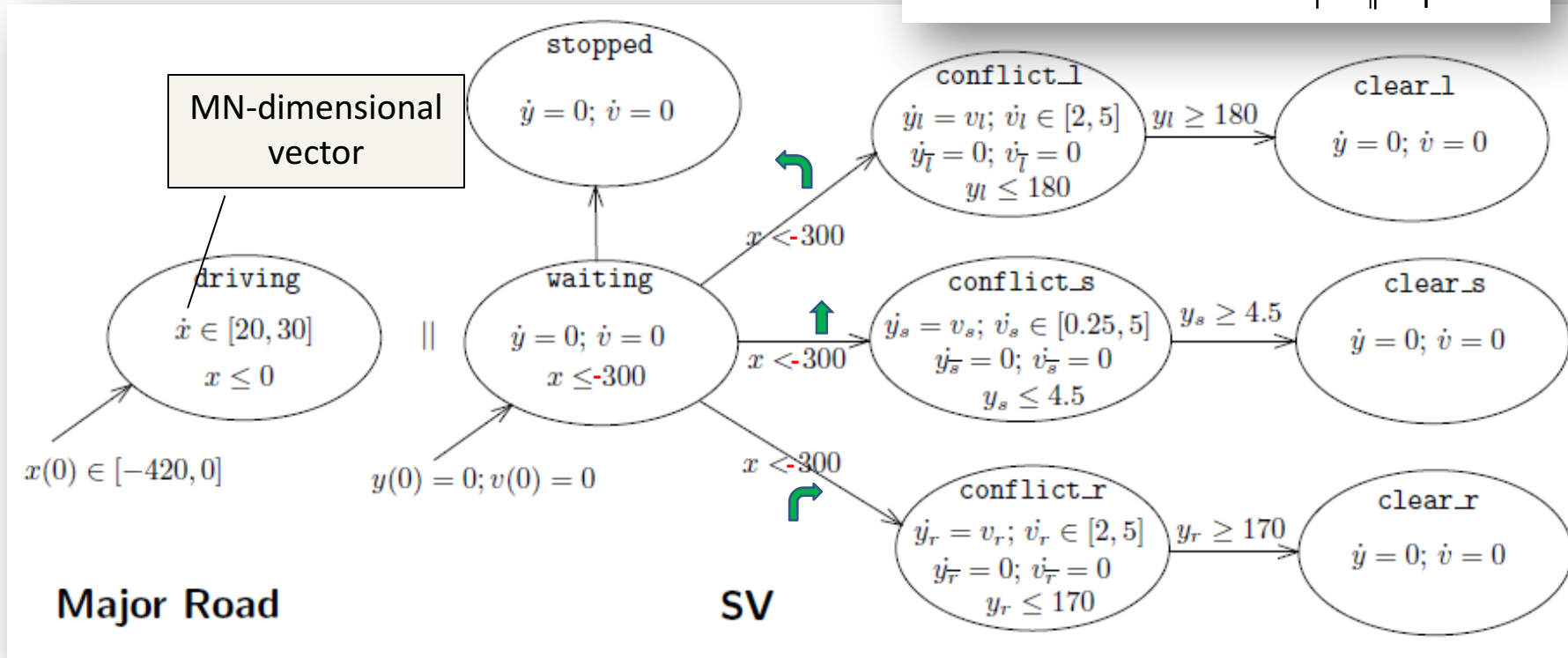
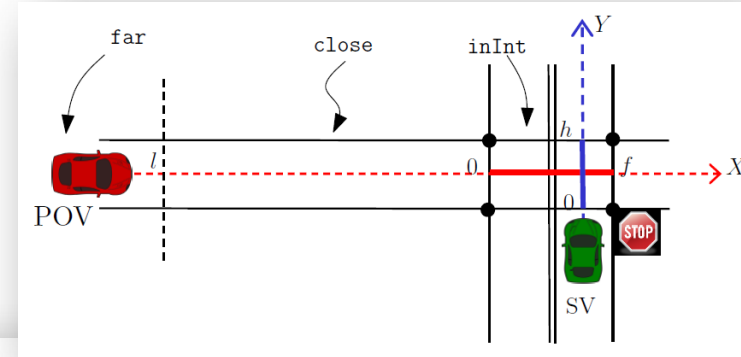
Hierarchical Verification

Conjunctive and disjunctive verification constructs can be nested arbitrarily





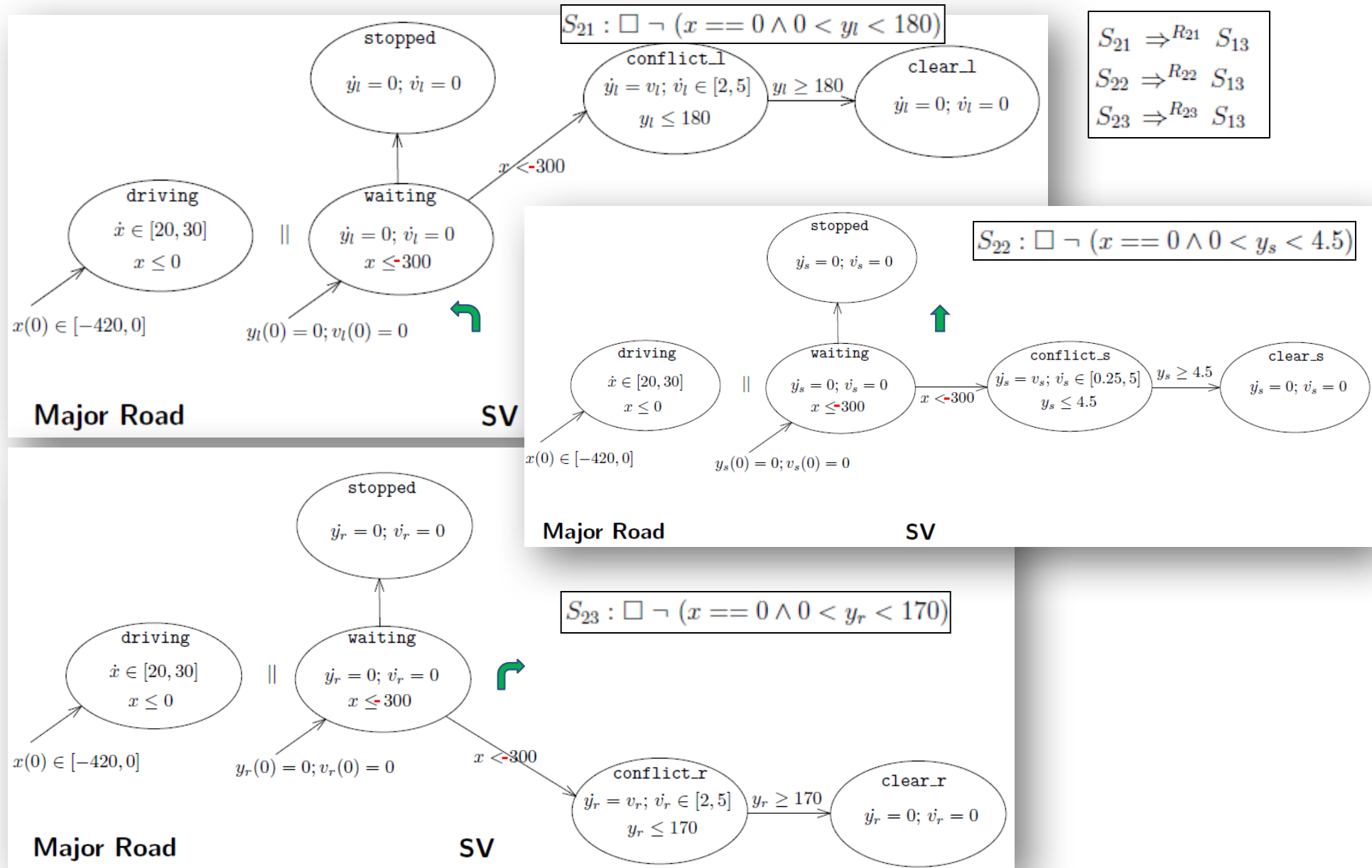
Verification Problem



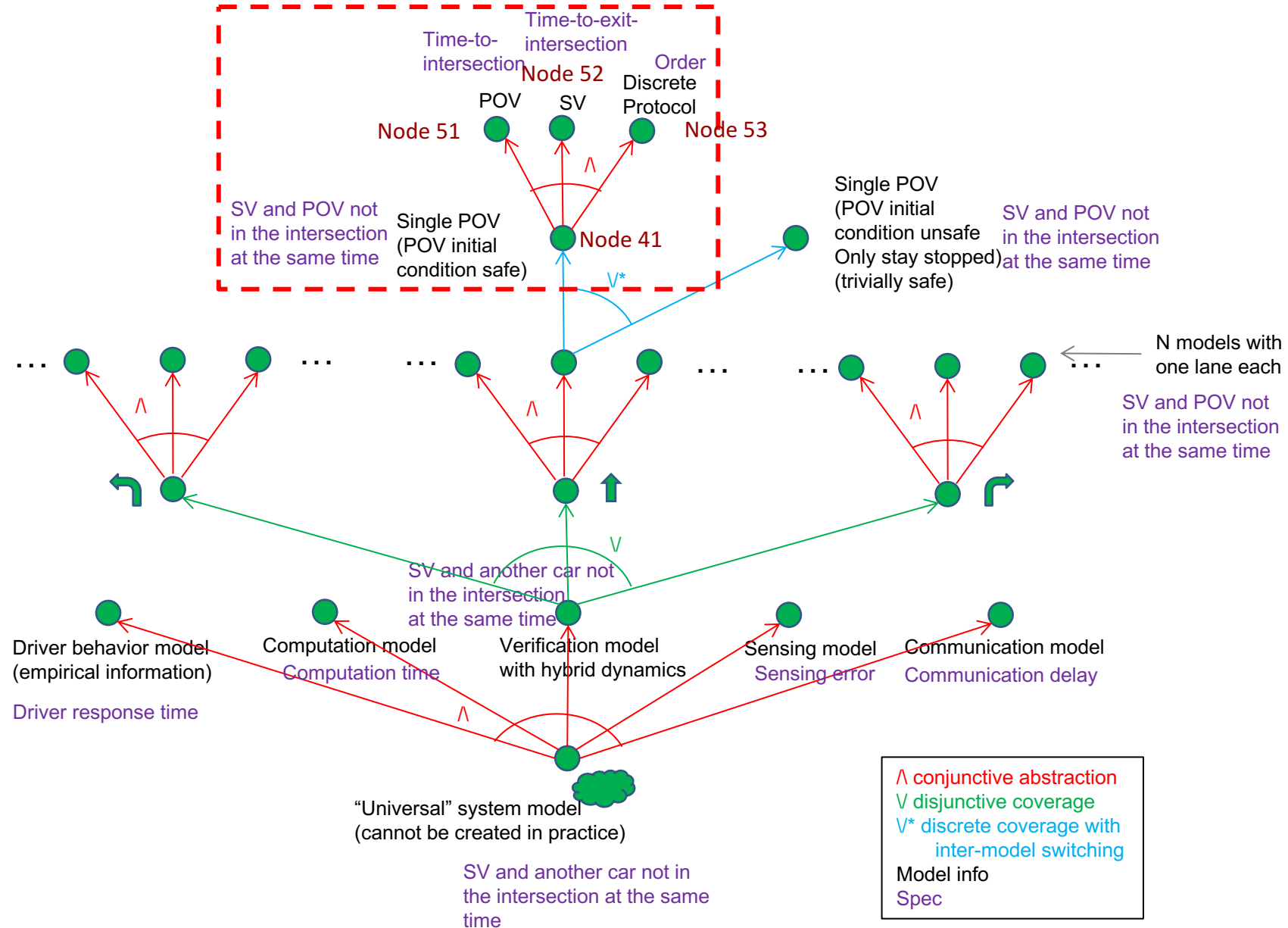
Verification objective: “SV and another car are never in the intersection at the same time”

$$\Box \neg ((x == 0 \wedge 0 < y_s < 4.5) \vee (x == 0 \wedge 0 < y_r < 170)) \vee (x == 0 \wedge 0 < y_l < 180))$$

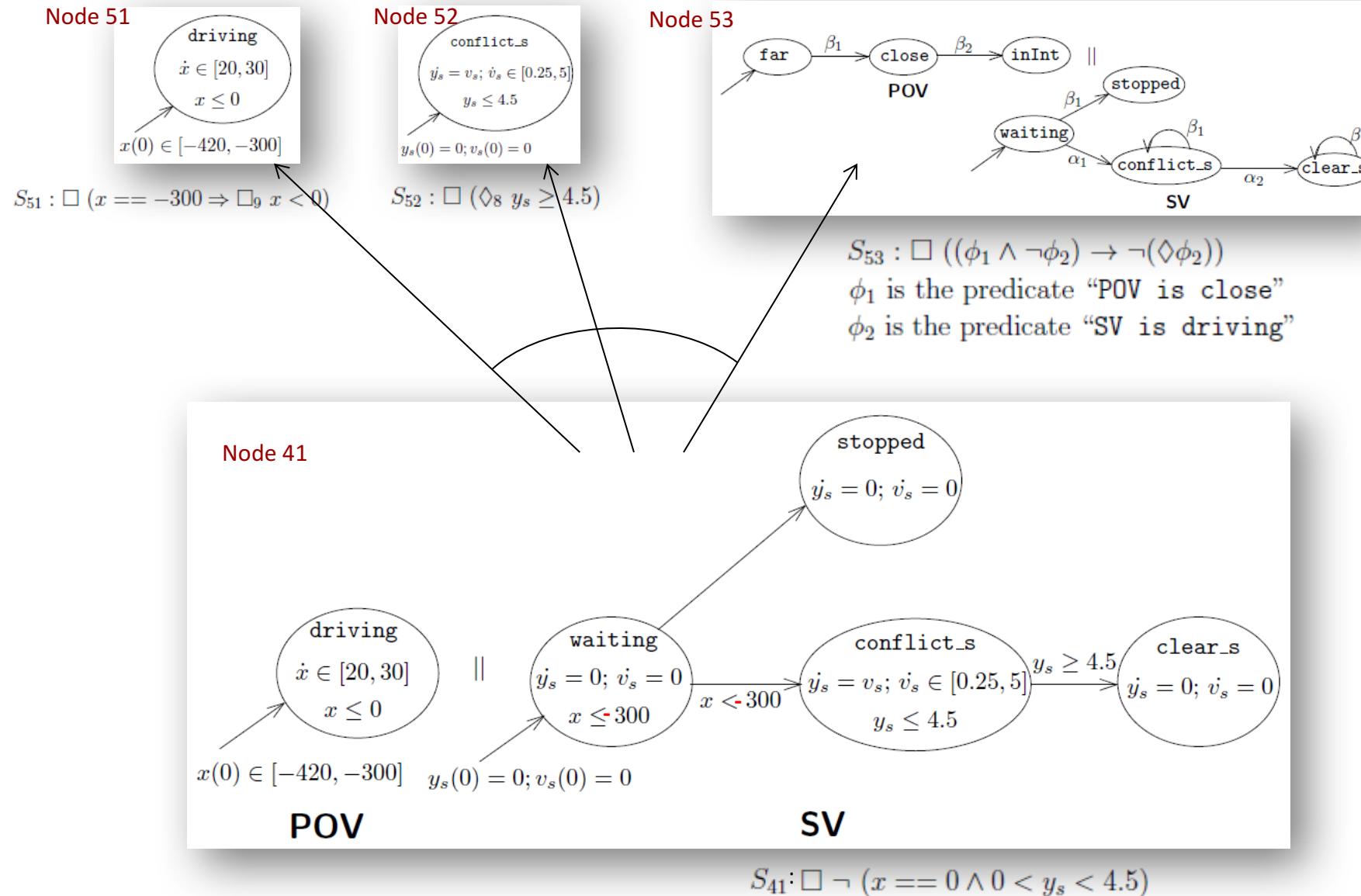
Disjunctive Heterogeneous Verification



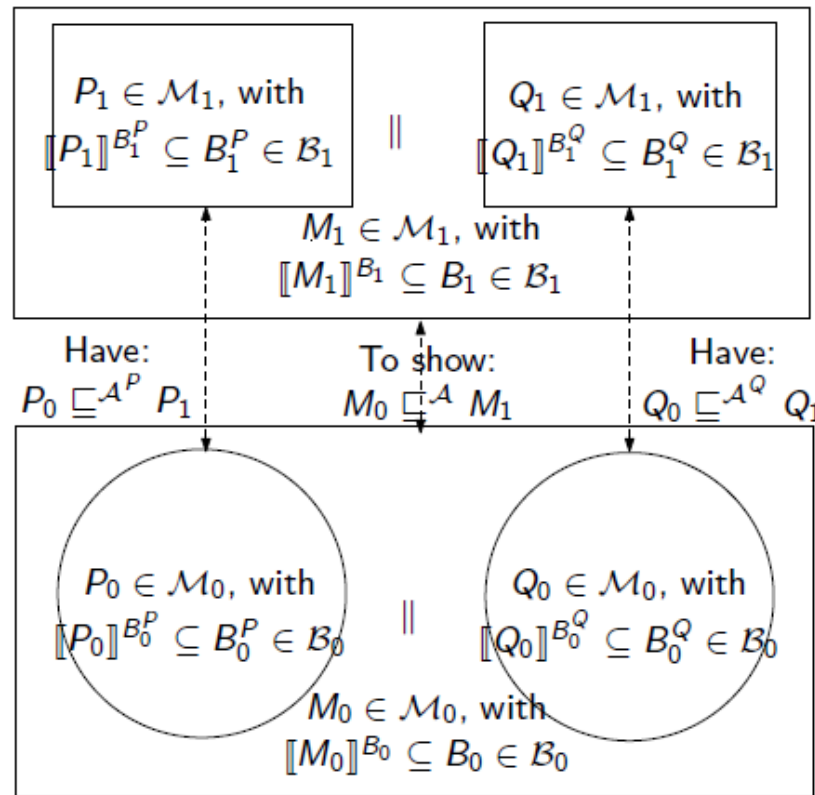
Heterogeneous verification of CICAS-SSA



Conjunctive Heterogeneous Verification



Leveraging Compositionality for Heterogeneous Abstraction



Schematic

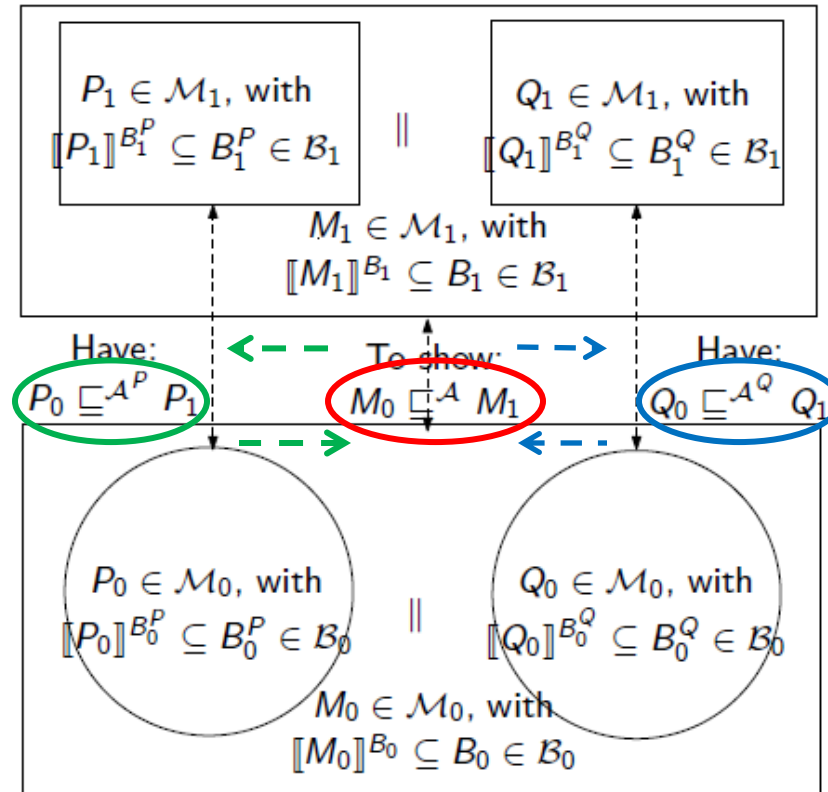
Objective: Conclude heterogeneous abstraction of the composition by establishing that of the components

Rationale: Component's local semantics defined in a behavior domain of smaller dimension

Need

- Behavior abstraction functions \mathcal{A} : behavior relations that are also functions
- Mappings between local/global behavior domains of the same type
- Mappings between local/global abstraction functions

Compositionality Conditions



★ “Models as composition of components”
 Analysis: Compositional Abstraction

Centralized Development

Start with \mathcal{A} , *localize* to get $\mathcal{A}^P, \mathcal{A}^Q$

If localizations of \mathcal{A} are \mathcal{A}^P and \mathcal{A}^Q , then compositional heterogeneous abstraction via \mathcal{A} holds

Decentralized Development

Start with $\mathcal{A}^P, \mathcal{A}^Q$, *globalize* to get \mathcal{A}

If globalizations of $\mathcal{A}^P, \mathcal{A}^Q$ are consistent (call it \mathcal{A}), then compositional heterogeneous abstraction via \mathcal{A} holds

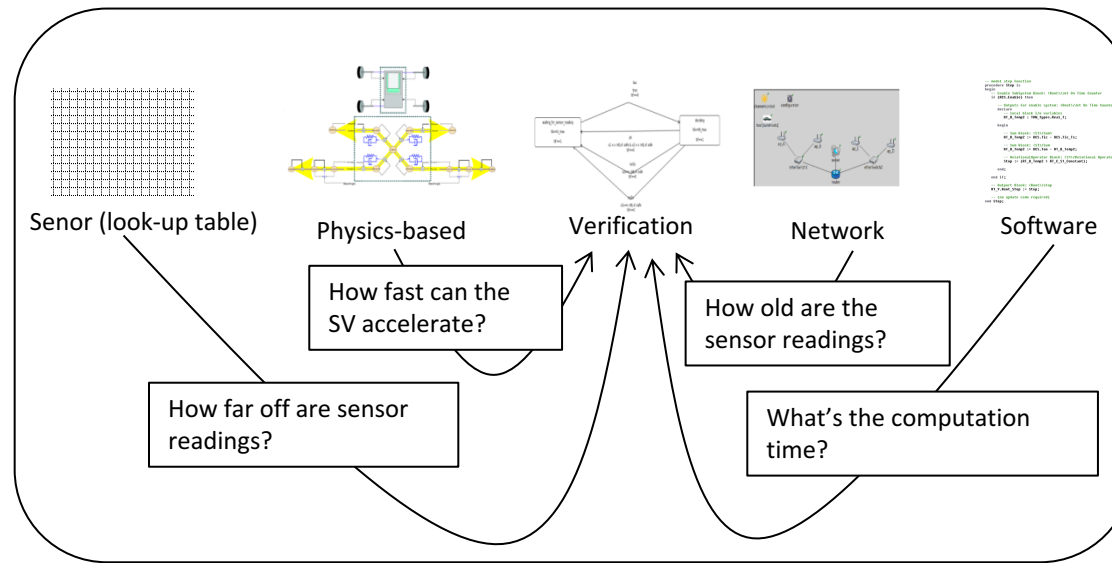
Compositional Heterogeneous Abstraction

HSCC '13

Akshay Rajhans

Bruce H. Krogh

Semantic Assumptions as Parameter Constraints



Problem

- *Semantic interdependencies* across formalisms
- *Consistency*

Challenge

- *Formal representation* that is *universal* to all modeling formalisms

Approach

- interdependencies as an *auxiliary constraint on parameters*
 - Find *effective constraint* on given model/spec. parameters (existential quantification)
 - Use *SMT solvers* or *theorem provers* to prove consistency
- ★ Ensures semantic (parameter) consistency using external SMT solvers or provers

Dependencies that cut across modeling formalisms can be captured as parameter constraints

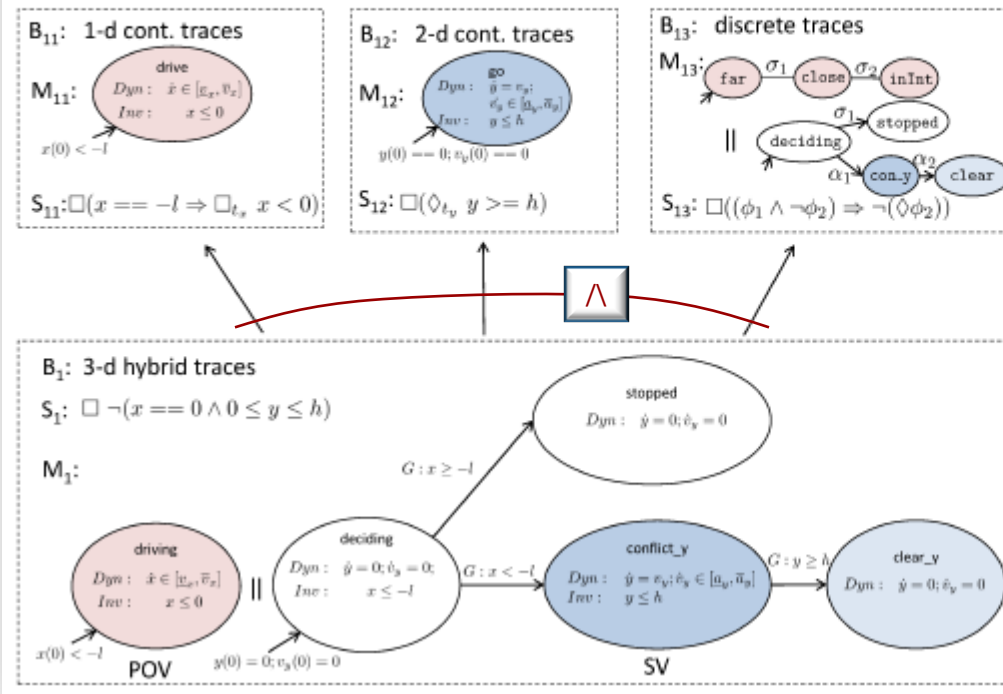
Using Parameters in Architectural Views to Support Heterogeneous Design and Verification

CDC '11

Akshay Rajhans[†], Ajinkya Bhawe[†], Sarah Loos[‡], Bruce H. Krogh[†], André Platzer[‡], David Garlan[‡]

Parametric Verification of CICAS

Parameterized models and specifications



1. Explicitly identify model parameters
 e.g. *speed limits, intersection geometry, minimum acceleration*, and spec. parameters, e.g., *POV min. time-to-intersection, SV max. time-to-clear-intersection*

2. Model interdependencies as an auxiliary constraint
 e.g., those dictated by *speed limits, newton's laws* and *intersection geometry* on *time-to-intersection*, ...

3. Project global constraints and interdependencies (aux. constraint) onto local sets of parameters

Heterogeneous Verification of Cyber-Physical Systems
 using Behavior Relations

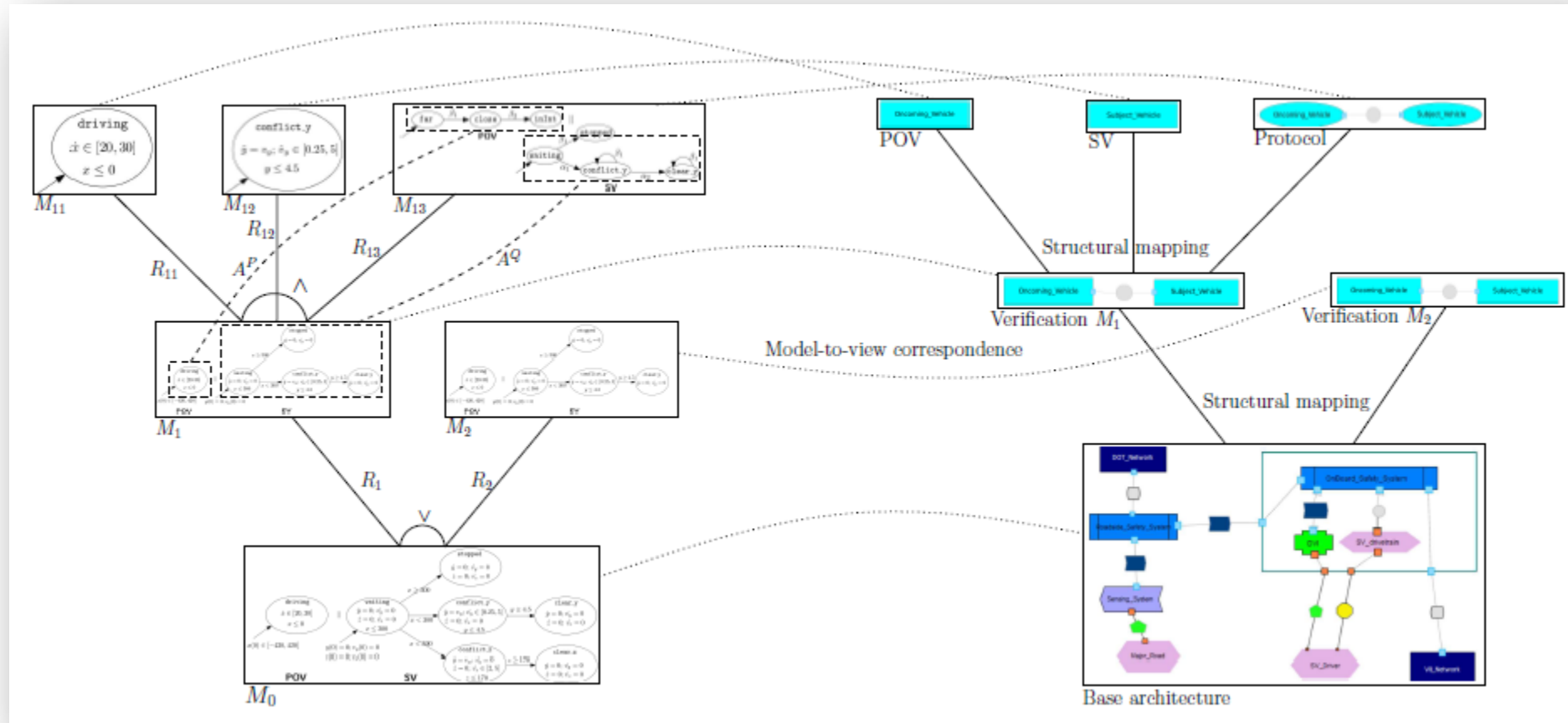
HSCC '12

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★ Proved semantic consistency in theorem prover KeYmaera

Semantic and Structural Hierarchies



Semantic side

Structural side

TAC '14
(CPS Special Issue)

Supporting Heterogeneity in Cyber-Physical Systems Architectures

Akshay Rajhans[†], Ajinkya Bhawe[†], Ivan Ruchkin[†], Bruce H. Krogh^{†*}, David Garlan[†], André Platzer[†] and Bradley Schmerl[†]

Summary

Cyber-Physical Systems present a major paradigm shift with systems that are

- Adaptive, Autonomous, Connected, and Collaborative

Model-based design critical for safe and efficient design process


- Open-ness and heterogeneity pose research challenges

Contributions for supporting heterogeneity in MBD of CPS

- *Architectural modeling*: high-level structural representation [MPM '09]
- *Model structures as architectural views* for comparing structure [ERTS '10]
- Semantic mappings using *behavior relations* enable *(compositional) heterogeneous verification* [HSCC '12, HSCC '13]
- *Constraint consistency* for consistent simplifying assumptions [CDC '11, HSCC '12]

Many challenges still remain

References*

 ECEASST

MPM '09

An Architectural Approach to the Design and Analysis of Cyber-Physical Systems

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Augmenting Software Architectures with Physical Components

Ajinkya Bhawe¹, David Garlan², Bruce H. Krogh¹, Akshay Rajhans¹, Bradley Schmerl²

ERTS² '10

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View Consistency in Architectures for Cyber-Physical Systems

ICCPs '11

Ajinkya Bhawe, Bruce H. Krogh David Garlan, Bradley Schmerl

Using Parameters in Architectural Views to Support Heterogeneous Design and Verification

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Heterogeneous Verification of Cyber-Physical Systems using Behavior Relations

HSCC '12

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Compositional Heterogeneous Abstraction

HSCC '13

Akshay Rajhans Bruce H. Krogh

Supporting Heterogeneity in Cyber-Physical Systems Architectures

TAC:CPS '13 (submitted)

Akshay Rajhans[†], Ajinkya Bhawe[†], Ivan Ruchkin[‡], Bruce H. Krogh^{†*}, David Garlan[‡], André Platzer[‡] and Bradley Schmerl[‡]

*Other work available at <https://arajhans.github.io>

